## 9<sup>TH</sup> TUTORIAL ON RANDOMIZED ALGORITHMS

Coupling of Markov chains

**1**. Random walk on a hypercube. In the hypercube graph of dimension d, vertices are binary strings of length d and two vertices are connected by an edge iff they differ in exactly one coordinate. For d = 2, the graph is

$$(\{00, 01, 10, 11\}, \{\{00, 01\}, \{00, 10\}, \{11, 01\}, \{11, 10\}\})$$

We start at  $0^d$  and do the following random walk:

- With probability 1/2 we stay at the current vertex.
- With probability 1/2 we choose an index  $j \in [d]$  uniformly at random and flip the *j*-th bit.

The Markov chain is ergodic (finite, aperiodic, and irreducible), so it converges to a unique stationary distribution (which one is it?). Show that the random walk has

$$\tau(\varepsilon) \le d \ln(d/\varepsilon)$$
.

**2**. Sampling colorings. Let G = (V, E) be a graph with maximum degree  $\Delta$ . We would like to uniformly sample a proper coloring with c colors. We use the following Markov chain: Given a proper coloring, pick  $v \in V$  and a color  $\ell \leq c$  uniformly at random and color v using  $\ell$  if this new coloring is proper.

- a) Show that this Markov chain is ergodic for  $c \ge \Delta + 2$ . What is the stationary distribution? What breaks down for  $c \le \Delta + 1$ ?
- b) Analyze the mixing time for  $c \ge 4\Delta + 1$  using a simple coupling (pick the same v and  $\ell$  in both chains). Hint: look at random variable  $d_t$  = the number of vertices with different colors in the two chains.
- c) Finally, we will show a better coupling that works even for  $c\geq 2\Delta+1$  (thus saving a factor of two).