## 4<sup>TH</sup> TUTORIAL ON RANDOMIZED ALGORITHMS

Exponentially decreasing bounds on the tails with Chernoff

**1**. Simulating a biased coin using a fair coin. We are given a fair coin with  $\Pr[\text{tails}] = 0.5$ . Show how to generate a random bit with  $\Pr[1] = p$  for a given  $p \in (0, 1)$  (both p = 0 and p = 1 are a bit boring). How many fair coin flips do we need in expectation? (This could be useful for some homework...)

- i) First, assume that  $p = k/2^{\ell}$  for some integers k and  $\ell$ .
- ii) What about any rational p, say, p = 1/3? What about irrational p?

**2**. Anti-Chernoff. Give an example of *n* dependent random variables  $X_i \in \{0, 1\}, i = 1, \ldots, n$  such that each  $X_i$  is a (fair or biased) coin flip but Chernoff bounds do not hold for  $\sum_i X_i$ . Are  $X_i$ 's positively or negatively correlated or uncorrelated?

**3**. Distinguishing coins. You are given two coins. One is fair and the other one has Pr[tails] = 1/4. We use the following algorithm to distinguish those:

- Pick a coin and toss it n times.
- Let  $\hat{p}$  be the probability of getting a tails (number of tails over n).
- If  $\hat{p} \ge 3/8$  we say this coin is fair.

Show that if  $n \ge 32 \ln(2/\delta)$  then our algorithm answers correctly with probability at least  $1 - \delta$ .

4. Estimator Theorem. Let U be a finite set and  $G \subseteq U$  its subset. We know |U| and wish to estimate |G|. If we take n uniformly random and independent samples from U (with replacement), calculate X = number of samples inside of G, and output  $A = X \frac{|U|}{n}$ . How large n do we need to choose so that A is a  $(1 \pm \varepsilon)$ -approximation of |G| with probability at least  $1 - \delta$ , i.e.,

$$\Pr\left[(1-\varepsilon)|G| \le A \le (1+\varepsilon)|G|\right] \ge 1-\delta?$$

5. QuickSort analysis. Consider the QuickSort algorithm with uniformly random pivot choice. Show that the expected number of comparisons is  $c \cdot n \ln(n)$  for some constant c. Prove that the probability of it making at least  $32n \ln(n)$  comparisons is at most  $1/n^3$ .