

# Probabilistic techniques - tutorials

## Classwork 6 – Markov chains

**Definition 1.** Let  $r_{i,j}^t = \Pr[X_t = j \text{ and } X_s \neq j \text{ for all } 0 < s < t | X_0 = i]$ . A state  $i$  is *recurrent* if  $\sum_{t \geq 1} r_{i,i}^t = 1$  and it is *transient* otherwise.

**Definition 2.** A communicating class  $C$  is *closed* if for all  $i \in C$  it holds that if  $j$  is accessible from  $i$ , then  $j$  is in  $C$  as well.

**Definition 3.** A communicating class  $C$  is *recurrent* if it contains a recurrent state. HW problem: If  $i \in C$  is recurrent, then every  $j \in C$  is recurrent.

1. Let  $\{X_i\}_{i=0}^\infty$  be a (homogeneous) Markov chain with the transition matrix  $P$ . Prove the following.

- (a) A state  $i$  is transient if and only if  $\sum_{n \geq 0} P_{i,i}^n < \infty$ . That is,  $i$  is recurrent if and only if  $\sum_{n \geq 0} P_{i,i}^n = \infty$ .

*Solution:*

- i. Let  $p_i = \Pr[\text{There exists } t \geq 1 \text{ such that } X_t = i | X_0 = i] = \sum_{t \geq 1} r_{i,i}^t$ .
- ii. Let  $N_i$  be the random variable that denotes the number of visits to  $i$ .
- iii. Compute expectation of  $N_i$ :  $\Pr[N_i = n | X_0 = i] = p_i^{n-1}(1 - p_i)$ , then

$$E[N_i | X_0 = i] = \sum_{n \geq 1} n p_i^{n-1} (1 - p_i) = (1 - p_i) \frac{d}{dp} \sum_{n \geq 1} p_i^n = \frac{1}{1 - p_i}.$$

- iv. We can compute the same expectation with indicator random variables:

$$E[N_i | X_0 = i] = E\left[\sum_{n \geq 1} I_{X_n=i} | X_0 = i\right] = \sum_{n \geq 1} \Pr[X_n = i | X_0 = i] = \sum_{n \geq 1} P_{i,i}^n.$$

- (b) A state  $i$  is recurrent if and only if

$$\Pr[X_n = i \text{ for infinitely many } n | X_0 = i] = 1.$$

Conversely, the state  $i$  is transient if and only if

$$\Pr[X_n = i \text{ for infinitely many } n | X_0 = i] = 0.$$

*Solution:*

- i.  $B_k = \{X_n = i \text{ for at least } k \text{ different values of } n \in \mathbb{N}\}$
- ii. Then  $\Pr[B_k | X_0 = i] = p_i^k$ .
- iii.  $B_1 \supset B_2 \supset B_3 \supset \dots$
- iv.  $\Pr[X_n = i \text{ for infinitely many } n | X_0 = i] = \Pr[\lim_k B_k | X_0 = i] = \lim_k \Pr[B_k | X_0] = \lim_k p_i^k$ .
- v. This last one is 1 iff  $p_i^k = 1$  if  $i$  is recurrent, and 0 otherwise.

2. Let  $i$  be a recurrent state and assume that  $j$  is accessible from  $i$ , then  $i$  is accessible from  $j$ . In particular recurrent communicating classes are closed.

*Solution:*

- (a) Since  $j$  is accessible from  $i$ , we can pick smallest  $k$  such that such that  $P_{i,j}^k > 0$ .
- (b) If  $i$  is not accessible from  $j$ , then  $p_i = \Pr[\text{returning to } i | X_0 = i] \leq 1 - P_{i,j}^k < 1$ .

(c) This is a contradiction with  $i$  being recurrent.

3. Every finite closed communicating class is recurrent.

*Solution:*

- (a) Let  $C$  be a closed communicating class with finitely many elements.
- (b) A chain starting in  $i$  stays in  $C$  forever and since  $C$  is finite, there must be at least one state  $j \in C$  which is visited infinitely often with positive probability, i.e.  $\Pr[X_n = j \text{ for infinitely many } n] > 0$ .
- (c) Since  $i$  and  $j$  are in the same communicating class, there exists  $m \in \mathbb{N}$  so that  $P_{j,i}^m > 0$ .
- (d) From the inequality

$$\Pr[X_n = j \text{ for infinitely many } n | X_0 = j] \geq P_{j,i}^m \Pr[X_n = j \text{ for infinitely many } n | X_0 = i] > 0$$

it follows that state  $j$  is recurrent by Exercise 1(b).

- (e) The class  $C$  is then recurrent because it contains at least one recurrent state, namely  $j$ .