

2.

Clique \Leftrightarrow IS

• a clique in G is an IS in G

3SAT \Leftrightarrow IS

- idea:
- if a formula is satisfiable, then in every clause there is a satisfied literal
 - for each clause we add a group of vertices
 - the independent set will pick a vertex from each group
 - if IS has size # clauses \Rightarrow each clause has a vertex of IS
 - \Rightarrow this IS vertex satisfies the clause

for clause $C_i = (l_{i1}, l_{i2}, l_{i3})$ we add three vertices u_{i1}, u_{i2}, u_{i3}
 and add edges $u_{i1}u_{i2}, u_{i2}u_{i3}, u_{i3}u_{i1}$

for two vertices $u_{i_a} \in C_i, u_{j_b} \in C_j$ ($i, j \in [3]$) we add an edge if they represent opposing literals

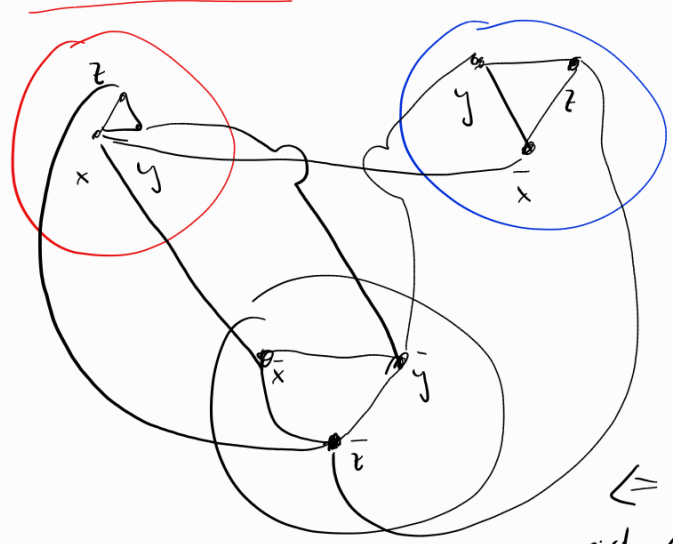
example.

$$Q = (x \vee y \vee \bar{z}) \wedge (\bar{x} \vee y \vee z) \wedge (\bar{x} \vee \bar{y} \vee \bar{z})$$

Claim: $m = \#$ clauses of Q .

Q is satisfiable \Leftrightarrow our graph has IS of size m

\Rightarrow If there is a sat assign, we can pick for each clause the vertex which corresponds to the satisfied literal.
 Sat assign has for each var either $x=0$ or $x=1 \Rightarrow$ there is no edge between two picked vertices



\Leftarrow If there is an IS of size m , it can pick ≤ 1 vertex from each clause. These vertices form a valid variable assign as there are no edges $x\bar{x}$.

IS \Rightarrow VC

the complement of IS in G is a VC

3.

let x_1, \dots, x_n be variables

if φ is not satisfiable \Rightarrow no, done

for $i \in [n]$:

if φ w/ $x_i = 1$ is sat:

set $x_i = 1$

else if φ w/ $x_i = 0$ is sat:

set $x_i = 0$

else:

this cannot happen because φ is sat so the sat assign either has

$x_i = 0$ or $x_i = 1$

4.

it is an FPT alg w/ param $n-k$

param k ? FPT alg probably doesn't exist (you will see later)

5.

suppose we recurse on u w/ neighbors u_1, \dots, u_d

we can bound the number on a n vertex gr. by

$$T(n) \leq 1 + T(n - \deg(u) - 1) + \sum_{i=1}^d T(n - \deg(u_i) - 1)$$

we choose a vertex u w/ min deg $\Rightarrow \deg(u_i) \geq \deg(u)$

set $s = \deg(u) + 1$

$$\rightarrow T(n) \leq 1 + s \cdot T(n-s) \leq 1 + s \cdot s^2 + \dots + s^{n/s} = O^*(s^{n/s})$$

• maximized when for integers $s=3 \rightarrow O^*(3^{n/3})$ w/g.