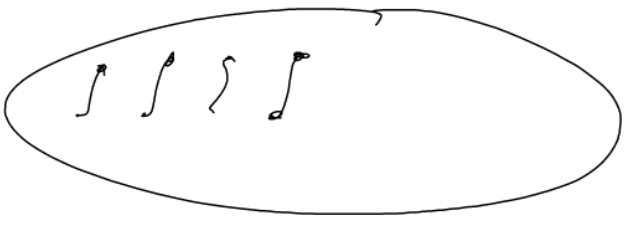


Packing & covering

• pack vertex disjoint copies of K_2



matching

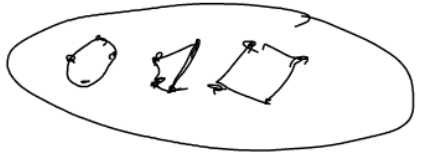
- Hall theorem
- Tutte theorem

cond. for general graphs

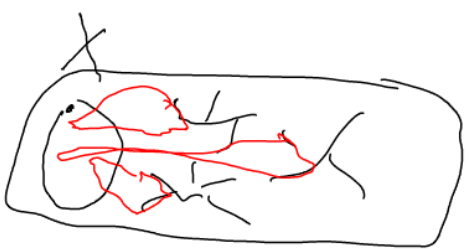


• packing of vertex disjoint cycles

Erdős-Pósa thm



$\exists f: \mathbb{N} \rightarrow \mathbb{N}$ s.t. $\forall G \neq G$ either G has k vertex disj. cycles
 or $\exists X \subseteq V(G) : G-X$ has no cycle
 & $|X| \leq f(k)$



G # of cycles in $G \leq |X|$

(not true with $f(k)=k$, unfortunately exercise)

• pack edge disj. paths $\Leftrightarrow k$ -edge connectivity (Menger's thm)



$|A| - |B| = k$

- pack vertex disj. paths $\Leftrightarrow k$ -vertex cover.
- packing dir. paths
max. flow = min cut



Packing & covering by spann. trees

① How many edge-disjoint spanning trees does a graph G have? MAX DST

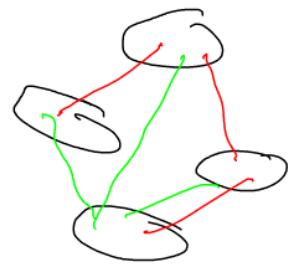
② How many sp. trees do we need to cover $E(G)$? MIN STC

Motivation $\left\{ \begin{array}{l} \text{nice } \& \ddot{\ } \\ k \text{ disj. sp. trees} \Rightarrow \text{thm 15} \\ k \text{ edge disj. u-o paths} \\ \& \text{ eff. algo} \end{array} \right.$

k edge disj. sp. t.s. $\Rightarrow k$ -edge cover.

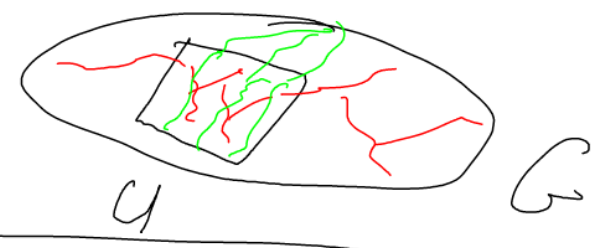
k -DST $\Rightarrow \forall$ part'n $V(G) = V_1 \cup \dots \cup V_r$
 \forall sp. tree has $\geq r-1$ edges u -between



$\Rightarrow \|G\| \geq k(r-1)$
 $|E(G)|$ G has $\geq k(r-1)$ edges u -between.

Thm 1 (Nash-Williams, Tutte 1961)
 Multigraph G has k -DST $\Leftrightarrow \forall$ part'n into r parts $\geq k(r-1)$ edges u -between
 conn.

k -STC $\Rightarrow \forall U \subseteq V(G) \quad \|G[U]\| \leq k(|U|-1)$



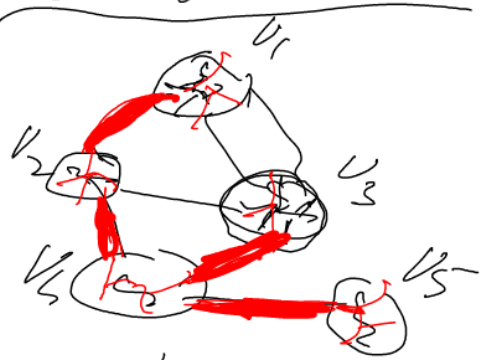
Thm 2 (Nash-Williams 1964)
 Multigraph G has STC by k trees $\Leftrightarrow \forall U \subseteq V(G) \quad \|G[U]\| \leq k(|U|-1)$

(Def arboricity of $G = \min k$ s.t.
 $E(G)$ can be covered by k sp-trees)

[Bowler, Carmesin, 2015]

Thm 3 $\forall G$ conn. multigraph $\forall k \exists$ partition \mathcal{P} of $V(G)$ s.t.

- ① $\forall u \in \mathcal{P}$ $G[u]$ has k edge disj. sp. trees
- ② G/\mathcal{P} has k -STC.

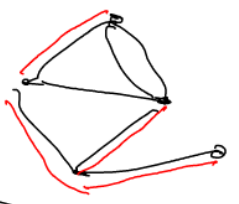


Thm 3 \Rightarrow Thm 2

G multigraph, conn.

\mathcal{P} for Thm 3

$\mathcal{P} = \{v_1, \dots, v_r\}$



Assume $\forall u : \|G[u]\| \leq k(|u|-1)$ edges

$u = v_i \Rightarrow G[u]$ has k edge disj. sp. trees. $\Rightarrow \|G[u]\| = k(|u|-1)$
 $G[u]$ is a union of k e.d.s.t.

② \Rightarrow sp. trees in G/\mathcal{P} T_1, \dots, T_k

T_1^i, \dots, T_k^i

$T_1^i = T_1 \cup T_1^i \cup \dots \cup T_1^r$ is a sp. tree in G

T_2^i
 \vdots
 T_k^i

cover of $E(G)$
 (edges of G/\mathcal{P} ... by Thm 3(2)
 edges in v_i

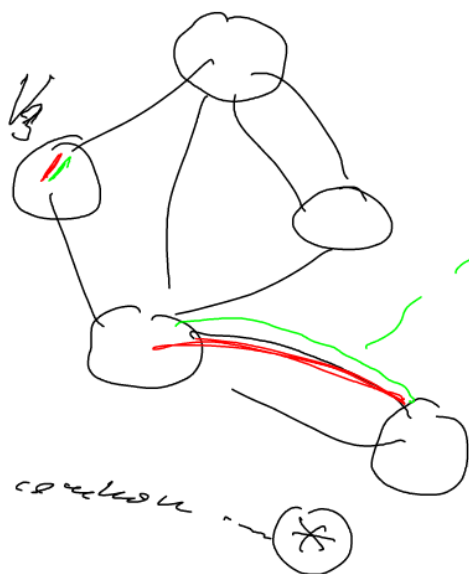
Thm 3 \Rightarrow Thm 1

$G, \mathcal{P} \rightarrow G/\mathcal{P}$ $\left\{ \begin{array}{l} \text{by Thm 3(2)} \dots T_1, \dots, T_k \text{ } k\text{-STC} \\ \text{by ass. of Thm 1} : \|G/\mathcal{P}\| \geq k(r-1) \end{array} \right\} \Rightarrow \|G/\mathcal{P}\| = k(r-1)$
 T_1, \dots, T_k are k -DST in G/\mathcal{P}

$G[V_i]$ has T_1^i, \dots, T_k^i k -DST \leftarrow $(*)$

$T_1 = T_1 \cup T_1^1 \cup \dots \cup T_1^r$ is a sp. tree in G

$T_2^1 \dots T_k^1$ ----- disjoint sp. trees



$T_1 \cap T_2^1$

has no edge in G/P

$T_1^3 \& T_2^3$

have no edge in common *

T --- sp. tree of G e --- chord ($e \notin E(T)$)

C_e --- fund. cycle of e w.r.t. T (the cycle of $T+e$)

$f \in C_e$ and $T' = T + e - f$ another sp. tree

$\mathcal{J} = (T_1, \dots, T_k) \longrightarrow E(\mathcal{J}) = \cup \{E(T_i) : T_i \in \mathcal{J}\}$

Seq. $e_0, e_1, \dots, e_n \in E(G)$ exchange chosen for \mathcal{J} started by e_0

- 1) $\forall i < n$ $\exists j = j(i)$ s.t. $e_i \in T_j$, e_{i+1} is a chord of T_j , its fund. cycle w.r.t. T_j contains e_i
- 2) $\{e_n \notin E(T_1) \cup \dots \cup E(T_k)\}$



$e_0 \in T_j$

e_1

$j(0) = 1$

$T_2^1 := T_2^0 + e_1 - e_0$

Lemma If e_0 starts an exchange chain for \mathcal{T}
 & lies in two of its trees
 then \exists family \mathcal{T}' of sp. trees s.t. $E(\mathcal{T}) \not\subseteq E(\mathcal{T}')$

- choose each chain e_0, e_1, \dots, e_n of max. length
 (given e_0)

- no e_i lies on pred. cycle w/out. any tree in \mathcal{T} of any
 e_j , $j > i+1$ (otherwise: $e_0, \dots, e_i, e_{i+1}, \dots, e_n$
 a shorter ex. chain)

- $\mathcal{T}^0 = \mathcal{T}$; \mathcal{T}^{i+1} ($i = 0, \dots, n-1$)

$$\begin{cases} \mathcal{T}^i = (T_1^i, \dots, T_k^i) & T_j^{i+1} = T_j^i + e_{i+1} - e_i \text{ for } j \neq i \\ \mathcal{T}^{i+1} = (T_1^{i+1}, \dots, T_k^{i+1}) & T_j^{i+1} = T_j^i \text{ otherwise} \end{cases}$$

- e_0 is in two trees \Rightarrow "delete" from one, repl. by e_1
 (now e_1 is in two trees)

----- $E(\mathcal{T}^n) = E(\mathcal{T}^0) \cup \{e_n\}$

(Proof of Thm 3) $\mathcal{T} = (T_1, \dots, T_k)$ family of k sp. trees of G ,
 $E(\mathcal{T})$ is maximal

$D = \{e \in E(G) : e \text{ starts an exch. chain for } \mathcal{T}\}$
 $\forall e \notin E(\mathcal{T}) : e \in D$ s.t. with $u=0$ ←

P -- part. of $V(G)$ into components of $(U(G), D)$

(1) $U \in P \quad S_j =$ subgraph of $T_j[U] \cap D$

$S_1 \dots S_k$ -- forests

-- edge disj.: (if $e \in S_1, S_2, \dots$ then $e \in D$)

$e = e$ starts an ex. ch.

by Lemma -- be part large T')

are S_i connected?

$\rightarrow E(D)$ form a conn. subgraph of $G[U]$

\rightarrow suffices to show $u, u' \in D, u, u' \in U \rightarrow \exists u-u'$ path in S_j .

1) $u, u' \in T_j \Rightarrow u, u' \in S_j$

2) $u, u' \in T_j$ path in T_j has all

its edges in D \Rightarrow all in S_j is

$\rightarrow e_0 = uu' \in D \rightarrow$ ex. ch. chain e_0, e_1, \dots, e_n

we want: $e \in D$

e, e_0, e_1, \dots, e_n is an ex. ch. chain

(2) recall: $U \in P \quad T_j[U]$ is a tree (S_j)

contracting U to a vertex, T_j stays a tree

$T_1 \dots T_k$ / P -- k-tuple of trees in G/P

want: $T'_1 \dots T'_k$ cover G/P

$e \in E \setminus E(D) \Rightarrow e \in D \Rightarrow \exists U \in P: e$ has both ends in $U \Rightarrow G/P$

Thm 3 [Bowler, Carmeson, 2015] $\#G$ conn. multigraph $\#k \exists$ partition \mathcal{P} of $U(G)$ st.

e is contracted

(1) $\forall U \in P \quad G[U]$ has k edge disj. sp. trees

(2) G/P has k -STC.

