

Our goal for this lecture is to present a (very brief) outline of the proof of the structure theorem. First, we need to formulate its local version with respect to a given tangle.

Let  $G$  be a graph and let  $\Omega$  be a cyclic sequence of vertices of  $G$ . Then  $(G, \Omega)$  is a *society*. We view a graph itself as a society with empty sequence. A *cross* in a society consists of two disjoint paths  $P_1$  and  $P_2$  in  $G$  such that the labels of the ends  $x_1$  and  $y_1$  of  $P_1$  and  $x_2$  and  $y_2$  of  $P_2$  can be chosen so that they appear in  $\Omega$  in order  $x_1, x_2, y_1$ , and  $y_2$ . A society is a *cell* if  $|\Omega| \leq 3$ . A *transaction* of order  $p$  in the society  $(G, \Omega)$  is a set of  $p$  pairwise vertex-disjoint paths with ends in  $\Omega$ . A society is a *p-vortex* if it contains no transaction of order greater than  $p$ ; from the homework assignment, we have the following description of *p*-vortices (recall the *adhesion* of a tree decomposition  $(T, \beta)$  is the maximum of  $|\beta(x) \cap \beta(y)|$  over distinct  $x, y \in V(T)$ ).

**Lemma 1.** *If  $(G, \Omega)$  is a  $p$ -vortex and  $\Omega = (v_1, \dots, v_m)$ , then  $G$  has a path decomposition  $(P, \beta)$  over the path  $P = v_1 v_2 \dots v_m$  of adhesion at most  $p$  such that  $v_i \in \beta(v_i)$  for  $i = 1, \dots, m$ .*

A society  $(G_1, \Omega_1)$  is a *subsociety* of  $(G, \Omega)$  if  $G_1$  is a subgraph of  $G$ , every edge of  $G_1$  incident with  $V(G_1) \setminus \Omega_1$  belongs to  $G_1$ , and  $G_1 \cap \Omega \subseteq \Omega_1$ . Two subsocieties  $(G_1, \Omega_1)$  and  $(G_2, \Omega_2)$  are *disjoint* if  $G_1 \cap G_2 = \Omega_1 \cap \Omega_2$ . A *segregation* of  $(G, \Omega)$  is a set  $\{(G_i, \Omega_i) : i = 1, \dots, n\}$  of its disjoint subsocieties such that  $G = G_1 \cup \dots \cup G_n$ . The segregation is *of type  $(k, p)$*  if all but at most  $k$  elements are cells and the remaining at most  $k$  elements are *p*-vortices.

If  $\Omega = \emptyset$ , an *arrangement* of the segregation in a surface  $\Sigma$  is a function  $\alpha$  satisfying the following conditions:  $\alpha(G_i, \Omega_i)$  is a disk  $\Delta_i \subseteq \Sigma$  and for each  $v \in \Omega_i$ ,  $\alpha(v)$  is a distinct point in  $\Sigma$  contained in the boundary of  $\Delta_i$ , such that

- for each  $i$ , the order of the points  $\alpha(v)$  for  $v \in \Omega_i$  in the boundary of  $\Delta_i$  matches the order of the vertices  $v$  in  $\Omega_i$ , and
- for distinct  $i$  and  $j$ , the disks  $\Delta_i$  and  $\Delta_j$  intersect exactly in the points  $\alpha(v)$  for  $v \in \Omega_1 \cap \Omega_2$ .

If  $\Omega$  is not emptyset, we additionally require  $\Sigma$  has exactly one hole and

- for each  $v \in \Omega$ , the point  $\alpha(v)$  is contained in the boundary of  $\Sigma$  and their order in the boundary matches the order of the vertices  $v$  in  $\Omega$ .

A society is *rural* if it has a segregation into cells with an arrangement in a disk. In the homework assignment, we have seen the following result.

**Lemma 2.** *A society  $(G, \Omega)$  is rural if and only if it does not contain a cross.*

For a tangle  $\mathcal{T}$  in  $G$  of order  $\theta$ , we say that a segregation  $\{(G_i, \Omega_i) : i = 1, \dots, n\}$  of  $G$  is  $\mathcal{T}$ -central if there is no  $(A, B) \in \mathcal{T}$  and  $i \in \{1, \dots, n\}$  such that  $B \subseteq G_i$ . For  $Z \subseteq V(G)$  with  $|Z| < \theta$ , recall that we can naturally define a tangle  $\mathcal{T} - Z$  in  $G - Z$  of order  $\theta - |Z|$  as the set of all separations  $\{(A - Z, B - Z) : (A, B) \in \mathcal{T}, Z \subseteq V(A \cap B)\}$ .

**Theorem 3** (The Structure Theorem, local version). *For every graph  $F$ , there exist integers  $\alpha < \theta$ ,  $k$ , and  $p$  such that the following holds. For every graph  $G$  and a tangle  $\mathcal{T}$  in  $G$  of order at least  $\theta$ , if  $F \not\subseteq G$ , then there exists  $A \subseteq V(G)$  of size at most  $\alpha$ , a surface  $\Sigma$  in which  $F$  cannot be drawn, and a  $(\mathcal{T} - A)$ -central segregation of  $G - A$  of type  $(k, p)$  with an arrangement in  $\Sigma$ .*

## 1 Global structure theorem from the local one

A graph  $G$  is  $(b, k, \rho)$ -near-embedded in a surface  $\Sigma$  if for some subset  $B \subseteq V(G)$ , the graph  $G - B$  has a drawing in  $\Sigma$  with at most  $k$  vortices of width at most  $\rho$ . The final global form of the structure theorem we aim for is as follows.

**Theorem 4** (The Structure Theorem, local version). *For every graph  $F$ , there exist integers  $b$ ,  $k$ , and  $\rho$  such that the following holds. For every graph  $G$ , if  $F \not\subseteq G$ , then  $G$  has a tree decompositions whose torsos can be  $(b, k, \rho)$ -near-embedded in surfaces in which  $F$  cannot be drawn. Equivalently,  $G$  can be obtained from graphs  $(b, k, \rho)$ -near-embedded in surfaces in which  $F$  cannot be drawn by clique-sums.*

In the first lecture, we have seen Theorem 4 follows from the following lemma, which as we now show is a consequence of Theorem 3. Recall a set  $\mathcal{L}$  of separations in a graph  $G$  is a *location* if for all distinct separations  $(A_1, B_1), (A_2, B_2) \in \mathcal{L}$ , we have  $A_1 \subseteq B_2$ . The *center* of the location is the graph  $C$  obtained from  $\bigcap_{(A,B) \in \mathcal{L}} B$  by adding all edges of cliques with vertex sets  $V(A \cap B)$  for  $(A, B) \in \mathcal{L}$ .

**Lemma 5.** *For every graph  $F$ , there exist integers  $\alpha < \theta$ ,  $k$ , and  $\rho$  such that the following holds. For every graph  $G$  and a tangle  $\mathcal{T}$  in  $G$  of order at least  $\theta$ , if  $F \not\subseteq G$ , then there exists a location  $\mathcal{L} \subseteq \mathcal{T}$  whose center is  $(\alpha, k, \rho)$ -near-embedded in a surface in which  $F$  cannot be drawn.*

*Proof.* Let  $\alpha < \theta$ ,  $k, p, A, \Sigma$ , and a  $(\mathcal{T} - A)$ -central segregation  $S$  of  $G - A$  of type  $(k, p)$  with an arrangement in  $\Sigma$  be obtained using Theorem 3. Let  $\rho = 2p + 1$ . The location  $\mathcal{L}$  is obtained as follows:

- For each cell  $(C, \Omega) \in S$ , we include the separation  $(A_C, B_C)$ , where  $A_C = G[V(C) \cup A]$  and  $V(A_C \cap B_C) = \Omega \cup A$ .
- For each  $p$ -vortex  $(C, \Omega) \in S$ , let  $(P, \beta)$  be the path decomposition from Lemma 1, where  $P = v_1 v_2 \dots v_m$ . In  $\mathcal{L}$ , we include all separations  $(A_i, B_i)$  for  $i = 1, \dots, m$ , where  $A_i = G[\beta(v_i) \cup A]$  and  $V(A_i \cap B_i) = A \cup \{v_i\} \cup X_i$ , with  $X_i = (\beta(v_i) \cap (\beta(v_{i-1}) \cup \beta(v_{i+1})))$ .

The  $(\alpha, k, \rho)$ -near-embedding of the center of  $\mathcal{L}$  is obtained by making  $A$  into apex vertices, replacing each cell in the arrangement by a clique of size at most three, and replacing each  $p$ -vortex by a vortex of width at most  $\rho$ , whose bags are the sets  $X_i$ .  $\square$

## 2 Growing animals

Let  $\mathcal{T}$  be a tangle in a graph  $G$ . For a surface  $\Sigma$ ,  $H \subseteq G$ , and a tangle  $\mathcal{T}_H$  of order  $\gamma$  in  $H$ , we say  $(H, \mathcal{T}_H)$  is a  $\Sigma$ -span of order  $\gamma$  in  $(G, \mathcal{T})$  if  $H$  is a subdivision of a 3-connected graph,  $H$  has a 2-cell drawing in  $\Sigma$  and  $\mathcal{T}_H$  is respectful for this drawing, and  $\mathcal{T}$  is conformal with  $\mathcal{T}_H$  (i.e., the tangle induced in  $G$  by  $\mathcal{T}_H$  is a subset of  $\mathcal{T}$ ). The results from the 6th lecture imply the following.

**Lemma 6.** *For every graph  $F$  and a surface  $\Sigma$  in which  $F$  can be drawn, there exists  $\gamma$  such that the following claim holds. If  $(G, \mathcal{T})$  contains a  $\Sigma$ -span of order  $\gamma$ , then  $F \preceq G$ .*

For a span  $(H, \mathcal{T}_H)$ , let  $d$  denote the distance function in  $H$  derived from  $\mathcal{T}_H$ . An  $H$ -path is a path in  $G$  intersecting  $H$  exactly in its endpoints. A  $(\gamma, s)$ -horn over the span is a vertex  $v \in V(G) \setminus V(H)$  for which there exist  $s$  paths from  $v$  to vertices  $v_1, \dots, v_s \in V(H)$ , disjoint except for their common start in  $v$  and disjoint from  $H$  except for their ends, where  $d(v_i, v_j) = \gamma$  for all  $i \neq j$ . For  $A \subseteq V(G) \setminus V(H)$ , a  $\gamma$ -hair avoiding  $A$  is a vertex  $z \in V(H)$  such that there exists an  $H$ -path in  $G - A$  to a vertex  $y \in V(H)$  with  $d(z, y) = \gamma$ .

A  $\Sigma$ -animal with  $a$  horns and  $b$  hairs of strength  $(\gamma, s)$  is a quadruple  $(H, \mathcal{T}_H, A, B)$ , where

- $(H, \mathcal{T}_H)$  is a  $\Sigma$ -span of order  $\gamma$ ,
- $A$  is a set of  $a$   $(\gamma, s)$ -horns over the span, and

- $B$  is a set of  $b$   $\gamma$ -hairs avoiding  $A$  such that  $d(x, y) = \sigma$  for distinct  $x, y \in B$ .

The argument used to prove Lemma 8 in the previous lecture notes gives the following.

**Lemma 7.** *For any surface  $\Sigma$  and integer  $m$ , there exist  $\gamma$  and  $s$  such that the following claim holds. If  $(G, \mathcal{T})$  contains a  $\Sigma$ -animal with  $\binom{m}{2}$  horns of strength  $(\gamma, s)$ , then  $K_m \preceq G$ .*

Furthermore, analogously to Lemma 9, we can show that many hairs can be combined into a horn.

**Lemma 8.** *For any surface  $\Sigma$  and integers  $m, s, \gamma, a$ , there exist  $\gamma'$  and  $b$  such that the following claim holds. If  $(G, \mathcal{T})$  contains a  $\Sigma$ -animal with  $a$  horns and  $b$  hairs of strength  $(\gamma', s + 1)$ , then either  $K_m \preceq G$ , or  $(G, \mathcal{T})$  contains a  $\Sigma$ -animal with  $a + 1$  horns of strength  $(\gamma, s)$ .*

Next, let us argue that  $H$ -paths between distant vertices of  $H$  either can be used to improve the animal, or can be all disrupted by a small number of vertices. We say a  $\Sigma$ -span  $(H, \mathcal{T}_H)$  is  $\delta$ -flat if for every  $H$ -path, its ends  $u$  and  $v$  satisfy  $d(u, v) < \delta$ . A  $\delta$ -zone around a vertex  $v \in V(H)$  is an open disk  $\Lambda \subset \Sigma$  bounded by a cycle  $C$  in  $H$  such that all atoms in the closure of  $\Lambda$  are at distance at most  $\delta$  from  $v$  and conversely, all atoms at distance at most  $\delta - 2$  belong to  $\Lambda$ . *Clearing the zone* means deleting vertices and edges of  $H$  drawn in  $\Lambda$ ; note that the resulting graph  $H'$  contains a tangle  $\mathcal{T}_{H'}$  of order  $\theta - O(\delta)$  conformal with  $\mathcal{T}_H$  such that the distances according to  $d_{\mathcal{T}_{H'}}$  are by at most  $O(\delta)$  smaller than those according to  $d_{\mathcal{T}_H}$ .

**Lemma 9.** *For any surface  $\Sigma$  and integers  $a, b, \theta_1$ , there exist  $\delta, \alpha$  and  $\theta_0$  such that the following claim holds for every integer  $s$ . If  $(G, \mathcal{T})$  contains a  $\Sigma$ -animal  $\mathcal{A} = (H, \mathcal{T}, A, B)$  with  $a$  horns and  $b$  hairs of strength  $(\theta, s)$  and  $\theta \geq \theta_0$ , then at least one of the following holds:*

1.  $(G, \mathcal{T})$  contains a  $(\Sigma + \text{handle})$ - or  $(\Sigma + \text{crosshandle})$ -span of order  $\theta_1$ ,  
or
2.  $(G, \mathcal{T})$  contains a  $\Sigma$ -animal with  $a$  horns and  $b + 1$  hairs of strength  $(\theta_1, s)$ , or
3.  $(G, \mathcal{T})$  contains a  $\Sigma$ -animal  $(H', \mathcal{T}_{H'}, A, B')$  with  $a$  horns and  $b$  hairs of strength  $(\theta - \delta, s)$  and a set  $Z \subseteq V(G) \setminus V(H')$  of size at most  $\alpha$  such that  $A \subseteq Z$  and  $(H', \mathcal{T}_{H'})$  is  $\delta$ -flat in  $G - Z$ .

*Proof idea.* For  $y \in B$ , choose a  $\delta'$ -zone  $\Lambda_y$  around  $y$  bounded by a cycle  $C_y$  for  $\delta' = O(\delta/b)$ , and let  $(H', \mathcal{T}_{H'})$  be the  $\Sigma$ -span of order  $\theta - \delta$  obtained by clearing the zones. For each  $y \in B$ , choose a vertex  $y'$  of  $C_y$  joined to  $y$  by a path in  $\Lambda_y$ , and let  $B' = \{y' : y \in B\}$ . Then  $(H', \mathcal{T}_{H'}, A, B')$  is an animal with  $a$  horns and  $b$  hairs of strength  $(\theta - \delta, s)$ . Moreover,  $\Lambda_y$  can be chosen so that  $C_y$  contains a set  $X_y$  that is free in  $\mathcal{T}_{H'}$  and satisfies the following technical connectivity condition ( $\star$ ): Suppose  $v_1, v_2, \dots, v_t$  is a sequence of vertices of  $H \cap \Lambda_y$  with  $d(y, v_1) \geq \delta - \theta_1$  and  $d(v_{j-1}, v_j) < \theta_1$  for  $j = 2, \dots, t$ . If  $d(y, v_m) < \theta_1$ , then  $H \cap (\Lambda_y \cup X_y)$  contains  $b^2\theta_1^2$  pairwise vertex-disjoint paths from  $\{v_1, \dots, v_m\}$  to  $X_y$ . Let us remark that to ensure this condition holds, it is in particular necessary that  $\delta \gg \theta_1$ .

If for some distinct  $y_1, y_2 \in B$  there are at least  $\theta_1^2$  disjoint  $H'$ -paths from  $X_{y_1}$  to  $X_{y_2}$ , then we can select  $\theta_1$  of them such that the order of their ends in  $C_{y_1}$  and  $C_{y_2}$  is either the same or opposite. Adding these paths to  $H'$ , we obtain a  $(\Sigma + \text{handle})$ - or  $(\Sigma + \text{crosshandle})$ -span of order  $\theta_1$  in  $G$ . Hence, we can assume this is not the case, and thus by Menger's theorem,  $G$  contains a set  $Z_0$  of size less than  $b^2\theta_1^2$  intersecting all paths with ends in  $X_{y_1}$  and  $X_{y_2}$  for distinct  $y_1, y_2 \in B$ .

If  $Z = (Z_0 \setminus V(H')) \cup A$  intersects all  $H'$ -path with ends  $u, v \in V(H')$  satisfying  $d_{\mathcal{T}_{H'}}(u, v) \geq \delta$ , then the last outcome holds. Hence, suppose  $Q$  is such a path avoiding  $Z$ . Let  $u'$  be the first vertex of  $Q$  after  $u$  belonging to  $H$ . If  $d(u, u') \geq \theta_1$ , we can add  $u$  as a hair to  $\mathcal{A}$  and the second outcome holds. Hence,  $d(u, u') < \theta_1$ , and defining  $v'$  symmetrically, we have  $d(v, v') < \theta_1$ . Consequently,  $d(u', v') > \delta - 2\theta_1 > 2\delta'$ . Therefore,  $u' \in \Lambda_{y_1}$  and  $v' \in \Lambda_{y_2}$  for distinct  $y_1, y_2 \in B$ . By the choice of  $Z_0$ , it cannot be the case that for both  $i \in \{1, 2\}$ ,  $H \cap (\Lambda_{y_i} \cup X_{y_i})$  contains a path from  $X_{y_i}$  to  $V(Q)$  disjoint from  $Z_0$ ; we can assume no such path exists for  $i = 1$ .

Consider the segment of  $Q$  starting with  $u'$  which intersects  $H$  only in  $H \cap \Lambda_{y_1}$ . Let  $v_1, \dots, v_m$  be these intersections in order they appear on  $Q$ , with  $m$  chosen maximum so that  $d(v_{j-1}, v_j) < \theta_1$  for  $j = 2, \dots, t$ . Since  $H \cap (\Lambda_{y_1} \cup X_{y_1})$  does not contain a path from  $X_{y_1}$  to  $V(Q)$  disjoint from  $Z_0$ , by ( $\star$ ) we have  $d(y_1, v_m) > \theta_1$ . Let  $w$  be the next intersection of  $Q$  with  $H$  after  $v_m$ . Then  $d(v_m, w) \geq \theta_1$ , by the maximality of  $m$  if  $w \in \Lambda_{y_1}$  and since  $d(y_1, y_3) \geq \theta$  and  $d(w_m, y_1) \leq \delta$  if  $w \in \Lambda_{y_3}$  for some  $y_3 \neq y_1$ . Then we can add  $v_m$  as a new hair to  $\mathcal{A}$ , and the second outcome holds.  $\square$

Let us now refine the last outcome. Suppose  $(H, \mathcal{T}_H)$  is a  $\Sigma$ -span. Another  $\Sigma$ -span  $(H', \mathcal{T}_{H'})$  is a  $\lambda$ -rearrangement of  $(H, \mathcal{T}_H)$  around a vertex  $v \in V(H)$  if every vertex or edge of  $H$  at distance more than  $\lambda$  from  $v$  belongs to  $H'$  and the distances according to  $d_{\mathcal{T}_{H'}}$  are by at most  $4\lambda + 2$  smaller than those according to  $d_{\mathcal{T}_H}$ . We say that  $(H, \mathcal{T}_H)$  is  $(\lambda, \delta)$ -flat if all its  $\lambda$ -rearrangements

are  $\delta$ -flat.

**Lemma 10.** *For any surface  $\Sigma$  and integers  $a, s, b, \theta_1$ , there exist  $\delta$  and  $\alpha$  such that for all  $\lambda$  and  $\theta_2$  there exists  $\theta$  for which the following claim holds. If  $(G, \mathcal{T})$  contains a  $\Sigma$ -animal  $\mathcal{A} = (H, \mathcal{T}, A, B)$  with  $a$  horns and  $b$  hairs of strength  $(\theta, s + \alpha)$ , then at least one of the following holds:*

1.  $(G, \mathcal{T})$  contains a  $(\Sigma + \text{handle})$ - or  $(\Sigma + \text{crosshandle})$ -span of order  $\theta_1$ ,  
or
2.  $(G, \mathcal{T})$  contains a  $\Sigma$ -animal with  $a$  horns and  $b + 1$  hairs of strength  $(\theta_1, s)$ , or
3. there is a set  $Z \subseteq V(G)$  of size at most  $\alpha$  such that  $(G - Z, \mathcal{T} - Z)$  contains a  $(\lambda, \delta)$ -flat  $\Sigma$ -span of order  $\theta_2$ .

*Proof idea.* Let  $\delta, \alpha$ , and  $\theta_0$  be as in Lemma 9, with  $\delta > \theta_1$ . Let  $\theta = \max(\theta_0, \theta_2 + \delta, 2\theta_1 + \delta + 4\lambda + 2)$ . Apply Lemma 9 to  $\mathcal{A}$ ; the first two outcomes correspond to the outcomes of this lemma, and thus we can assume  $(G, \mathcal{T})$  contains a  $\Sigma$ -animal  $(H', \mathcal{T}_{H'}, A, B')$  with  $a$  horns and  $b$  hairs of strength  $(\theta - \delta, s + \alpha)$  and a set  $Z \subseteq V(G) \setminus V(H')$  of size at most  $\alpha$  such that  $A \subseteq Z$  and  $(H', \mathcal{T}_{H'})$  is  $\delta$ -flat in  $G - Z$ . If the  $\Sigma$ -span  $(H', \mathcal{T}_{H'})$  is  $(\lambda, \delta)$ -flat in  $G - Z$ , then the third outcome holds.

Otherwise, there exists a  $\lambda$ -rearrangement  $(H'', \mathcal{T}_{H''})$  of  $(H', \mathcal{T}_{H'})$  around a vertex  $w$  and an  $H''$ -path  $Q$  in  $G - Z$  whose ends  $u$  and  $v$  satisfy  $d_{\mathcal{T}_{H''}}(u, v) \geq \delta$ . For a vertex  $x \in A$ , consider  $s + 1$  of the paths showing that  $x$  is a horn which are disjoint from  $Z \setminus \{x\}$ . At most one of them intersects  $H'' - H'$ , as otherwise  $(H', \mathcal{T}_{H'})$  would not be  $\delta$ -flat in  $G - Z$ . Consequently, each element of  $A$  is a  $(\theta_1, s)$ -horn over  $(H'', \mathcal{T}_{H''})$ . Furthermore, at most one of the hairs in  $B'$  is at distance less than  $\theta_1 + \lambda$  from  $w$ ; let  $y$  be this hair, or an arbitrary hair in  $B'$  if no hair is close to  $w$ . Then  $(H'', \mathcal{T}_{H''}, A, (B' \setminus \{y\}) \cup \{u, v\})$  is a  $\Sigma$ -animal with  $a$  horns and  $b + 1$  hairs of strength  $(\theta_1, s)$ , and the second outcome of the lemma holds.  $\square$

For a span  $(H, \mathcal{T}_H)$ , a face of  $H$  is an *eye* if there exist vertices  $x_1, \dots, x_4$  appearing in order in the cycle bounding  $f$  such that there exist disjoint  $H$ -paths from  $x_1$  to  $x_3$  and from  $x_2$  to  $x_4$  and the set  $\{x_1, \dots, x_4\}$  is free in  $\mathcal{T}_H$ . The homework assignment for Lesson 6 implies the following result (the assumption that the span is  $(\beta - 10)$ -flat is used to show that the crossing paths for different eyes are pairwise disjoint).

**Lemma 11.** *For any surface  $\Sigma$  and integer  $m$ , if  $\beta \gg m, \Sigma$  and  $(G, \mathcal{T})$  contains a  $(\beta - 10)$ -flat  $\Sigma$ -span with  $m^4$  eyes pairwise at distance at least  $\beta$  apart, then  $K_m \preceq \mathcal{T}$ .*

We now can improve or embed a flat span.

**Lemma 12.** *For any surface  $\Sigma$  and integers  $m, \delta$ , and  $\theta_1$ , there exist  $\lambda, \theta_2$ , and  $p$  such that the following claim holds. If  $(G, \mathcal{T})$  contains a  $(\lambda, \delta)$ -flat  $\Sigma$ -span of order at least  $\theta_2$  and  $K_m \not\leq G$ , then either*

1.  $(G, \mathcal{T})$  contains a  $(\Sigma + \text{crosscap})$ -span of order  $\theta_1$ , or
2.  $G$  has a  $\mathcal{T}$ -central segregation of type  $(m^4, p)$  with an arrangement in  $\Sigma$ .

*Proof idea.* Choose  $\gamma_{-1} \gg \lambda_0 \gg \gamma_0 \gg \lambda_1 \gg \gamma_1 \gg \dots \gg \lambda m^4 \gg \gamma_{m^4}$ , where  $\gamma_{m^4} \geq \max(\delta, \beta) + 10$  for  $\beta$  from Lemma 11. Set  $\theta_2 = \gamma_{-1}$  and  $p \gg \lambda = \lambda_0$ .

Let  $k \in \{0, \dots, m^4\}$  be maximum such that  $(G, \mathcal{T})$  contains a  $(\lambda_k, \delta)$ -flat  $\Sigma$ -span  $(H, \mathcal{T}_H)$  of order at least  $\gamma_{k-1}$  with  $k$  eyes  $f_1, \dots, f_k$  pairwise at distance at least  $\gamma_k$  apart. Note that  $k < m^4$ , as otherwise  $K_m \leq G$  by Lemma 11.

Consider a vertex  $v \in V(H)$  such that an eye can be created by cleaning a  $4\delta$ -zone around  $v$ . The maximality of  $k$  implies that the distance between  $v$  and some  $f_i$  is less than  $\gamma_{k+1} + 16\delta + 2 \ll \lambda_{k+1}$ . For  $i = 1, \dots, k$ , let  $\Lambda_i$  be the corresponding zone around  $f_i$ . The ‘‘local planarity’’ together with rigidness of  $H$  implies that everything outside of these zones can be broken up into cells with an arrangement in  $\Sigma$ . We now apply the results from the homework assignment to each zone  $\Lambda_i$  and all  $H$ -bridges of  $G$  that attach to it. If it does not contain a large crooked transaction, then it can be decomposed into a rural neighborhood and a  $p$ -vortex, and if this happens for all  $i$ , we obtain a desired  $\mathcal{T}$ -central segregation of type  $(m^4, p)$  with an arrangement in  $\Sigma$ .

Hence, suppose this does not happen for some  $i$ . The large crooked transaction contains one of crosscap, jump, or double-cross type. Crosscap-type crooked transaction can be used to rearrange  $(H, \mathcal{T}_H)$  into a  $(\Sigma + \text{crosscap})$ -span of order  $\theta_1$ . Jump one would contradict the assumption that  $(H, \mathcal{T}_H)$  is  $(\lambda_k, \delta)$ -flat. The double-cross type can be used to rearrange and obtain one more distant eye, contradicting the maximality of  $k$ .  $\square$

Let us now combine all these results.

**Corollary 13.** *For any surface  $\Sigma$  and integers  $m, \theta_1, s_1, a, b$ , there exists  $\theta, s, p, \alpha$  such that the following holds. If  $(G, \mathcal{T})$  contains a  $\Sigma$ -animal with  $a$  horns and  $b$  hairs of strength  $(\theta, s)$  and  $K_m \not\leq G$ , then*

- $(G, \mathcal{T})$  contains a  $(\Sigma + \text{handle})$ -,  $(\Sigma + \text{crosshandle})$ -, or  $(\Sigma + \text{crosscap})$ -span of order  $\theta_1$ , or

- $(G, \mathcal{T})$  contains a  $\Sigma$ -animal with  $a$  horns and  $b + 1$  hairs of strength  $(\theta_1, s_1)$ , or
- there exists a set  $Z \subseteq V(G)$  of size at most  $\alpha$  such that  $G - Z$  has a  $(\mathcal{T} - Z)$ -central segregation of type  $(m^4, p)$  with an arrangement in  $\Sigma$ .

*Proof.* Let  $\delta$  and  $\alpha$  be as in Lemma 10 for the given parameters (with  $s = s_1$ ). Let  $\lambda$ ,  $\theta_2$ , and  $p$  be as in Lemma 12 for the given parameters. Let  $\theta$  be obtained by Lemma 10 for this  $\lambda$  and  $\theta_2$ , and let  $s = s_1 + \alpha$ .

Applying Lemma 10 to  $(G, \mathcal{T})$ , the first outcomes correspond to the first two outcomes of this lemma. Hence, we can assume the third outcome, giving us a  $(\lambda, \delta)$ -level  $\Sigma$ -span in  $(G - Z, \mathcal{T} - Z)$  for a set  $Z$  of size at most  $\alpha$ . We now apply Lemma 12; the first outcome corresponds to the first outcome of this lemma, while the second one corresponds to the third outcome of this lemma.  $\square$

### 3 Proof of the structure theorem

Let  $m = |V(F)|$ . We define  $\theta = \theta(\Sigma, a, b)$ ,  $s = s(\Sigma, a, b)$ ,  $p = p(\Sigma, a, b)$  and  $\alpha = \alpha(\Sigma, a, b)$  inductively so that the following conditions hold.

- If  $F$  can be drawn in  $\Sigma$ , then  $\theta$  is equal to  $\gamma$  from Lemma 6 and  $s = p = \alpha = 0$ . Suppose from now on that  $F$  cannot be drawn in  $\Sigma$ .
- If  $a \geq \binom{m}{2}$ , then let  $\theta$  and  $s$  be chosen according to Lemma 7, setting  $\theta = \gamma$ , and let  $p = \alpha = 0$ . Suppose from now on that  $a < \binom{m}{2}$ .
- Let  $b_{\max}(\Sigma, a)$  be equal to  $b$  from Lemma 8 for the given  $\Sigma$ ,  $m$ ,  $s(\Sigma, a + 1, 0)$ ,  $\theta(\Sigma, a + 1, 0)$ , and  $a$ . If  $b \geq b_{\max}(\Sigma, a)$ , then let  $s = s(\Sigma, a + 1, 0) + 1$ ,  $p = \alpha = 0$ , and let  $\theta$  be chosen as  $\gamma'$  from Lemma 8. From now on, suppose that  $b < b_{\max}(\Sigma, a)$ .
- Let  $\theta$  and  $s$  be chosen according to Corollary 13, with  $\theta_1$  and  $s_1$  maximum of the following:
  - $\theta(\Sigma', 0, 0)$  and  $s(\Sigma', 0, 0)$  for  $\Sigma' \in \{\Sigma + \text{handle}, \Sigma + \text{crosshandle}, \Sigma + \text{crosscap}\}$ .
  - $\theta(\Sigma, a, b + 1)$  and  $s(\Sigma, a, b + 1)$ .

We choose  $p$  and  $\alpha$  as the maximum of values of  $p$  and  $\alpha$  among these cases and those obtained from Corollary 13.

A straightforward inductive argument gives the following.



**Corollary 14.** *For any graph  $F$ , a surface  $\Sigma$ , and integers  $a$  and  $b$ , if  $(G, \mathcal{T})$  contains a  $\Sigma$ -animal with  $a$  horns and  $b$  hairs of strength  $(\theta(\Sigma, a, b), s(\Sigma, a, b))$  and  $F \not\subseteq G$ , then there exists a set  $Z \subseteq V(G)$  of size at most  $\alpha(\Sigma, a, b)$  such that  $G - Z$  has a  $\mathcal{T}$ -central segregation of type  $(m^4, p(\Sigma, a, b))$  with an arrangement in some surface in which  $F$  cannot be drawn.*

Theorem 3 thus follows with  $\alpha = \alpha(\mathbf{sphere}, 0, 0)$ ,  $k = |V(F)|^4$ ,  $p = p(\mathbf{sphere}, 0, 0)$ , and  $\theta$  large enough that (by the grid theorem), any graph with a tangle of order at least  $\theta$  contains a wall of order  $\theta(\mathbf{sphere}, 0, 0)$ —such a wall forms a **sphere**-animal with no horns and hairs.