

Observation

$$\omega(G) \leq \chi(G) \leq \Delta(G) + 1$$

Lemma

There exist triangle-free graphs ($\omega = 2$) with arbitrarily large chromatic number.

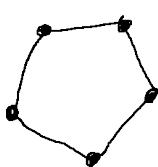
Which graphs satisfy $\chi(G) = \omega(G)$?

- Problem: $\chi(G + K_{|V(G)|}) = \omega(G + K_{|V(G)|}) = |V(G)|$

Definition

A graph G is **perfect** if every induced subgraph H of G satisfies $\chi(H) = \omega(H)$.

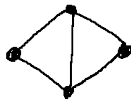
Q: Which of the following graphs are perfect?



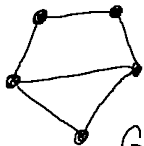
G_1



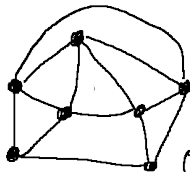
G_2



G_3



G_4



G_5

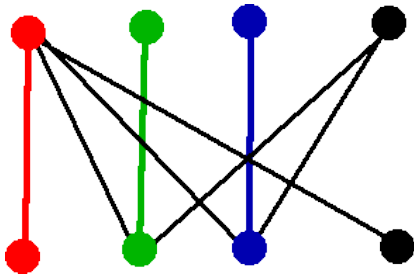
Examples: chordal graphs, bipartite graphs, complements of bipartite graphs, linegraphs of bipartite graphs, comparability graphs, ...

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$$\chi(G) = \omega(G) = 2$$

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$$\chi(\overline{G}) = |V(G)| - \beta(G) = \alpha(G) = \omega(\overline{G})$$

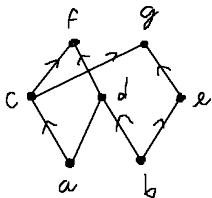


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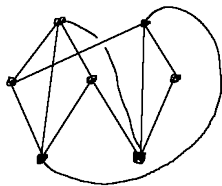
$$\chi(L(G)) = \chi'(G) = \Delta(G) = \omega(L(G))$$

Examples: chordal graphs, bipartite graphs, complements of bipartite graphs, linegraphs of bipartite graphs, comparability graphs, ...

Comparability graph: For a partial ordering \prec , $uv \in E(G)$ iff $u \prec v$ or $v \prec u$.



$$\begin{array}{l}
 a < c < f \quad b < d < c < g \\
 a < d < f \quad b < e < g
 \end{array}$$



G

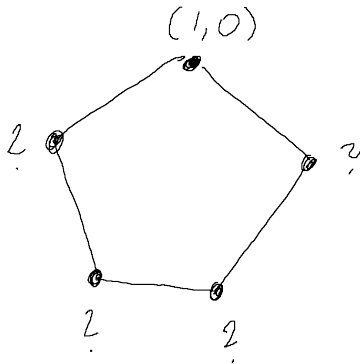
$\omega(G)$ = longest chain, $\chi(G)$ = partition to min. # of antichains.

Definition

For $r \geq 2$, a **vector r -coloring** is a function $\varphi : V(G) \rightarrow$ unit vectors such that for every $uv \in E(G)$,

$$\langle \varphi(u), \varphi(v) \rangle \leq -1/(r-1).$$

Vector $\sqrt{5}$ -coloring of C_5 ; $-1/(\sqrt{5}-1) = \cos \frac{4\pi}{5}$:



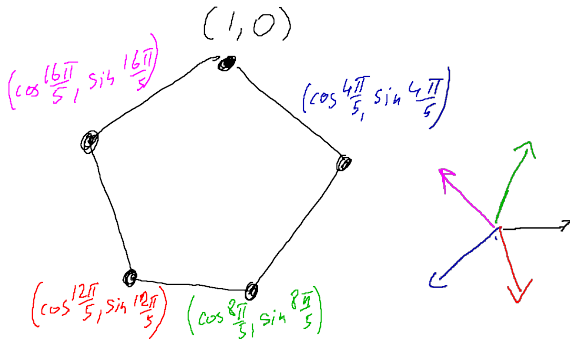
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Definition

Vector chromatic number

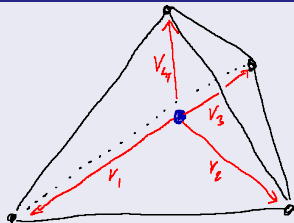
$$\chi_v(G) = \inf\{r \geq 2 : G \text{ has a vector } r\text{-coloring}\}.$$

Lemma

For G with $E(G) \neq \emptyset$, $\chi_v(G) \leq \chi(G)$.

Proof.

v_1, \dots, v_c : unit vectors forming a regular simplex, $\langle v_i, v_j \rangle = s$.



$$0 = \left| \sum_{i=1}^c v_i \right|^2 = \sum_{i=1}^c |v_i|^2 + \sum_{i \neq j} \langle v_i, v_j \rangle = c + c(c-1)s$$

$$s = -1/(c-1)$$

Vertices of color $i \rightarrow v_i$.



Lemma

$$\chi_v(G) \geq \omega(G)$$

Proof.

v_1, \dots, v_k : Vectors on a clique in a vector r -coloring.

$$\begin{aligned} 0 &\leq \left| \sum_{i=1}^k v_i \right|^2 = \sum_{i=1}^k |v_i|^2 + \sum_{i \neq j} \langle v_i, v_j \rangle \\ &\leq k - k(k-1)/(r-1) \\ k &\leq r \end{aligned}$$



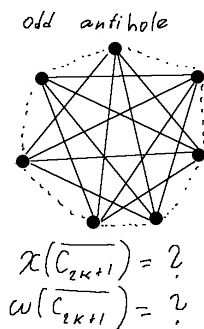
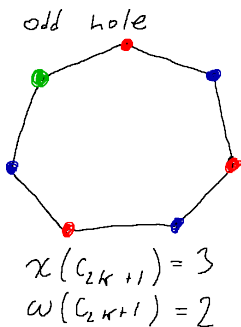
- $\omega(G) \leq \chi_v(G) \leq \chi(G)$
- $\omega(G) = \chi(G) \Rightarrow$ equal to $\chi_v(G)$.
- $\chi_v(G)$ can be computed by semidefinite programming.

$$\chi_v(G) = 1 - 1/t,$$

where t is minimum s.t. there exist vectors $\{v_z : z \in V(G)\}$ satisfying

$$\begin{array}{ll} \langle v_z, v_z \rangle = 1 & \text{for every } z \in V(G) \\ \langle v_y, v_z \rangle \leq t & \text{for every } yz \in E(G) \end{array}$$

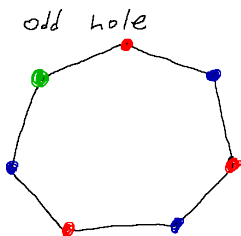
- **Hole**: induced (≥ 4) -cycle
- **Antihole**: induced subgraph isomorphic to the complement of a (≥ 4) -cycle



Theorem (Strong Perfect Graph Theorem)

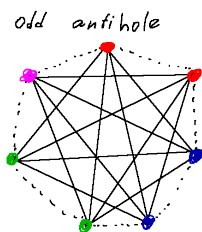
A graph is perfect if and only if it contains no odd holes and antiholes.

- **Hole**: induced (≥ 4) -cycle
- **Antihole**: induced subgraph isomorphic to the complement of a (≥ 4) -cycle



$$\chi(C_{2k+1}) = 3$$

$$\omega(C_{2k+1}) = 2$$



$$\chi(\overline{C_{2k+1}}) = k+1$$

$$\omega(\overline{C_{2k+1}}) = k$$

Theorem (Strong Perfect Graph Theorem)

A graph is perfect if and only if it contains no odd holes and antiholes.

Corollary (Weak Perfect Graph Theorem)

A graph is perfect if and only if its complement is perfect.

Lemma

$A_1, \dots, A_k, K_1, \dots, K_k \subseteq \{1, \dots, n\}$. If

- $A_i \cap K_j = \emptyset$ for $i \in \{1, \dots, k\}$ and
- $|A_i \cap K_j| = 1$ for $i, j \in \{1, \dots, k\}, i \neq j$,

then $k \leq n$.

Proof.

- WLOG $k \geq 2$.
- S $k \times n$ matrix, $S_{i,j} = 1$ if $j \in A_i$, 0 otherwise.
- T $n \times k$ matrix, $T_{i,j} = 1$ if $i \in K_j$, 0 otherwise.
- $(ST)_{i,j} = |A_i \cap K_j|$: 0 on the diagonal, 1 elsewhere.
- $k = \text{rank}(ST) \leq \text{rank}(S) \leq n$.



Lemma

G is perfect if and only if for every induced subgraph H of G , there exists an independent set intersecting all cliques of size $\omega(H)$.

Proof.

- \Rightarrow Clique of size $\omega(H)$ intersects all color classes of $\omega(H)$ -coloring.
- \Leftarrow
- Induction on $|V(G)|$.
 - A intersects all cliques of size $\omega(G)$:
 $\omega(G - A) = \omega(G) - 1$.
 - Induction hypothesis:
 $\chi(G - A) = \omega(G - A) = \omega(G) - 1$.
 - Color A using a new color: $\chi(G) = \omega(G)$.



Lemma

G is perfect if and only if for every induced subgraph H of G,

$$\alpha(H)\omega(H) \geq |V(H)|.$$

G perfect: $\omega(H) = \chi(H) \geq |V(H)|/\alpha(H)$.

Lemma

G is perfect if and only if for every induced subgraph H of G ,

$$\alpha(H)\omega(H) \geq |V(H)|.$$

- For every independent set A , there exists a disjoint $\omega(G)$ -clique $K(A)$
- $A_0 = \{v_1, \dots, v_\alpha\}$ an independent set.
- Induction: $\chi(G - v_i) \leq \omega$; color classes $A_{i,1}, \dots, A_{i,\omega}$.
 - $K(A_0)$ intersects all of them.
 - $K(A_{i,j})$ intersects all but $A_{i,j}$.
 - $K(A_{i',j})$ for $i' \neq i$ intersects all of them.
 - $K(A_{i',j}) \cap A_0 = \{v_{i'}\} \Rightarrow K(A_{i',j}) \subseteq V(G - v_i)$
- Lemma for sets A_\star and $K(A_\star)$: $\alpha\omega + 1 \leq |V(G)| \nmid$

Lemma

G is perfect if and only if for every induced subgraph H of G ,

$$\alpha(H)\omega(H) \geq |V(H)|.$$

If $\alpha(H)\omega(H) \geq |V(H)|$, then

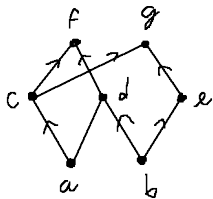
$$\alpha(\overline{H})\omega(\overline{H}) = \omega(H)\alpha(H) \geq |V(H)|.$$

Corollary (Weak Perfect Graph Theorem)

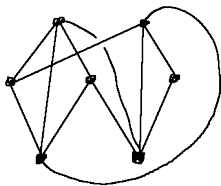
If G is perfect, then \overline{G} is perfect.

G : the comparability graph of \prec

- $\omega(\overline{G}) = \alpha(G) =$ largest antichain in \prec
 - The **width** of \prec .
- $\chi(\overline{G}) =$ partition of elements to smallest number of chains



$$\begin{array}{l} a < c < f \quad b < d \quad c < g \\ a < d < f \quad b < e < g \end{array}$$



G

Corollary (Dilworth's theorem)

Every finite partially ordered set of width k has a partition into k chains.