

# Coloring of triangle-free graphs and the Rosenfeld counting method

Zdeněk Dvořák

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There are many constructions that show that triangle-free graphs can have arbitrarily large chromatic number. Let us give one that shows that being triangle-free does not help to get better bound on the chromatic number even for degenerate graphs.

**Lemma 1.** *For every integer  $d \geq 1$ , there exists a  $d$ -degenerate graph  $G_d$  of girth at least six and of chromatic number  $d + 1$ .*

*Proof.* We prove the claim by induction on  $d$ . For  $d = 1$ , we can take  $G_1 = K_2$ . Suppose now that  $d \geq 2$ . The graph  $G_d$  is obtained as follows. We start with an independent set  $K$  of size  $d(|V(G_{d-1})| - 1) + 1$  and with  $\binom{|K|}{|V(G_{d-1})|}$  copies of  $G_{d-1}$ . For each subset  $S$  of  $K$  of size  $|V(G_{d-1})|$ , we then choose a distinct copy  $G_S$  of  $G_{d-1}$  and add a perfect matching between  $S$  and  $V(G_S)$ . Observe that  $G_d$  has girth at least six, as any cycle is either contained in a copy of  $G_{d-1}$ , or contains edges of at least two copies of  $G_{d-1}$  and passes through  $K$ . Moreover,  $G_d$  is  $d$ -degenerate, as each copy of  $G_{d-1}$  is  $(d - 1)$ -degenerate and the degrees of its vertices are increased only by 1 in  $G_d$ .

Suppose for a contradiction that  $G_d$  can be colored using at most  $d$  colors. By the pigeonhole principle, some set  $S \subseteq K$  of size  $|V(G_{d-1})|$  is colored using just one color. But then the copy  $G_S$  of  $G_{d-1}$  would be colored using only the remaining  $d - 1$  colors, which is a contradiction.  $\square$

However, we can get an improvement in terms of the maximum degree. We give a relatively recent proof of Martinsson, which uses an idea based on a counting argument of Rosenfeld. To illustrate the method, let us start with a simpler example. Recall that a *star coloring* is a proper coloring in which any two color classes induce a star forest (or equivalently, at least three colors appear on any 4-vertex path).

**Theorem 2.** *For any positive integer  $\Delta$ , any graph of maximum degree at most  $\Delta$  has a star coloring by at most  $\lceil 13\Delta^{3/2} \rceil$  colors.*

*Proof.* Let  $k = \lceil 13\Delta^{3/2} \rceil$  and  $\beta = k/3$ . Let  $\mathcal{C}(G)$  be the set of all star colorings of  $G$  using colors  $\{1, \dots, k\}$ . We will prove that if  $G$  is a graph of maximum degree at most  $\Delta$ , then for every  $v \in V(G)$ , we have

$$\frac{|\mathcal{C}(G)|}{|\mathcal{C}(G-v)|} \geq \beta.$$

This clearly implies  $\mathcal{C}(G) \neq \emptyset$ , and in fact that  $G$  has at least  $\beta^{|V(G)|}$  star colorings by at most  $k$  colors. We prove the claim by induction on the number of vertices of  $G$ . When  $|V(G)| = 1$ , we have  $|\mathcal{C}(G)| = k > \beta$  and  $|\mathcal{C}(G-v)| = 1$ , and thus the claim holds. Hence, we can assume that  $|V(G)| > 1$ .

Let  $\mathcal{C}_v(G)$  be the set of proper  $k$ -colorings of  $G$  whose restriction to  $G-v$  is a star coloring. Note that any coloring in  $\mathcal{C}(G-v)$  can be extended to a coloring in  $\mathcal{C}_v(G)$  by choosing a color of  $v$  different from the color of the neighbors of  $v$ , and this can be done in at least  $k - \Delta$  ways. Hence,

$$|\mathcal{C}_v(G)| \geq (k - \Delta)|\mathcal{C}(G-v)|.$$

For a 4-vertex path  $P$  containing  $v$ , let  $\mathcal{C}_P$  be the set of colorings in  $\mathcal{C}_v$  that use only two colors on  $P$ . Letting  $\mathcal{P}$  be the set of all 4-vertex paths in  $G$  containing  $v$ , we have

$$\mathcal{C}(G) = \mathcal{C}_v \setminus \bigcup_{P \in \mathcal{P}} \mathcal{C}_P.$$

Hence, we need to bound  $|\mathcal{C}_P|$ . Note that each coloring in  $\mathcal{C}_P$  is obtained from a star coloring of  $G - V(P)$  by choosing the two colors used on  $P$ , and thus (using the induction hypothesis),

$$|\mathcal{C}_P| \leq k^2 |\mathcal{C}(G - V(P))| \leq \frac{k^2}{\beta^3} |\mathcal{C}(G-v)|.$$

Putting these bounds together, we have

$$\begin{aligned} |\mathcal{C}(G)| &\geq |\mathcal{C}_v| - \sum_{P \in \mathcal{P}} |\mathcal{C}_P| \\ &\geq (k - \Delta)|\mathcal{C}(G-v)| - |\mathcal{P}| \cdot \frac{k^2}{\beta^3} |\mathcal{C}(G-v)| \\ &= \left( k - \Delta - 2\Delta^3 \cdot \frac{k^2}{\beta^3} \right) |\mathcal{C}(G-v)| \\ &\geq \beta |\mathcal{C}(G-v)|. \end{aligned}$$

□

Let us now give a (somewhat more involved) argument for the chromatic number of triangle-free graphs.

**Theorem 3.** *For every  $\Delta \geq 10^{10}$ , every triangle-free graph of maximum degree at most  $\Delta$  has chromatic number at most  $\lceil 4\Delta/\log \Delta \rceil$ .*

*Proof.* Let  $k = \lceil 4\Delta/\log \Delta \rceil$  and  $\ell = \frac{\Delta^{1-\log 2}}{\log \Delta}$ . For a graph  $G$ , let  $\mathcal{C}(G)$  be the set of all proper  $k$ -colorings of  $G$ . We will show that if  $G$  is a triangle-free graph of maximum degree at most  $\Delta$ , then for every  $v \in V(G)$ , we have

$$\frac{|\mathcal{C}(G)|}{|\mathcal{C}(G-v)|} \geq \ell.$$

This clearly implies  $\mathcal{C}(G) \neq \emptyset$ , and in fact that  $G$  has at least  $\ell^{|V(G)|}$   $k$ -colorings. We prove the claim by induction on the number of vertices of  $G$ . When  $|V(G)| = 1$ , we have  $|\mathcal{C}(G)| = k > \ell$  and  $|\mathcal{C}(G-v)| = 1$ , and thus the claim holds. Hence, we can assume that  $|V(G)| > 1$ .

For a vertex  $x \in V(G)$  and a partial coloring  $\varphi$  of  $G$ , let  $a(G, \varphi, x)$  denote the number of colors in  $\{1, \dots, k\}$  that do not appear on the neighbors of  $x$  in  $\varphi$ . Note that each coloring  $\varphi \in \mathcal{C}(G-v)$  extends to a proper  $k$ -coloring of  $G$  in exactly  $a(G, \varphi, v)$  ways, and thus

$$\frac{|\mathcal{C}(G)|}{|\mathcal{C}(G-v)|} = \frac{\sum_{\varphi \in \mathcal{C}(G-v)} a(G, \varphi, v)}{|\mathcal{C}(G-v)|} = E[a(G, \varphi, v)],$$

where the expectation is over a  $k$ -coloring  $\varphi$  of  $G-v$  chosen uniformly at random. Hence, we need to prove that

$$E[a(G, \varphi, v)] \geq \ell.$$

Let  $t = \frac{\ell}{\log \Delta} \geq 2$ ; we say a neighbor  $u$  of  $v$  is  $\varphi$ -poor if  $a(G-v, \varphi, u) \leq t$ . Consider any  $k$ -coloring  $\psi$  of  $G-v-u$ . If  $a(G-v, \psi, u) > t$ , then  $\psi$  does not extend to any  $k$ -coloring  $\varphi$  of  $G-v$  such that  $u$  is  $\varphi$ -poor; if  $a(G-v, \psi, u) \leq t$ , then  $\psi$  extends to exactly  $a(G-v, \psi, u) \leq t$   $k$ -colorings  $\varphi$  of  $G-v$  such that  $u$  is  $\varphi$ -poor. Hence, using the induction hypothesis for  $G-v$  and  $u$ , the number of  $k$ -colorings of  $G-v$  such that  $u$  is  $\varphi$ -poor is at most

$$t \cdot |\mathcal{C}(G-v-u)| \leq \frac{t}{\ell} |\mathcal{C}(G-v)|,$$

and thus

$$\Pr[u \text{ is } \varphi\text{-poor}] \leq \frac{t}{\ell}.$$

Let  $q(\varphi)$  denote the number of  $\varphi$ -poor neighbors of  $v$ . Then

$$E[q(\varphi)] \leq \frac{t}{\ell} \cdot \deg v \leq \frac{t\Delta}{\ell} \leq \frac{k}{4},$$

and thus

$$\Pr \left[ q(\varphi) > \frac{k}{2} \right] \leq \frac{1}{2}.$$

We say that  $v$  is  $\varphi$ -rich if  $v$  has at most  $\frac{k}{2}$   $\varphi$ -poor neighbors. Hence,

$$\Pr[v \text{ is } \varphi\text{-rich}] \geq \frac{1}{2}.$$

Consequently,

$$E[a(G, \varphi, v)] \geq \Pr[v \text{ is } \varphi\text{-rich}] \cdot E[a(G, \varphi, v) | v \text{ is } \varphi\text{-rich}] \geq \frac{1}{2} E[a(G, \varphi, v) | v \text{ is } \varphi\text{-rich}],$$

and thus it suffices to show that

$$E[a(G, \varphi, v) | v \text{ is } \varphi\text{-rich}] \geq 2\ell.$$

Note that since  $G$  is triangle-free, the neighborhood  $N(v)$  of  $v$  is an independent set, and thus whether  $v$  is  $\varphi$ -rich depends only on the restriction of  $\varphi$  to  $V(G) \setminus N[v]$ . Hence, it suffices to show that for every  $k$ -coloring  $\psi_0$  of  $G - N[v]$  such that  $v$  is  $\psi_0$ -rich and  $\psi_0$  can be extended to a  $k$ -coloring of  $G - v$ ,

$$E[a(G, \varphi, v) | \varphi \text{ extends } \psi_0] \geq 2\ell.$$

Let  $R$  be the set of the neighbors of  $v$  that are not  $\psi_0$ -poor. Observe it suffices to show that for any  $k$ -coloring  $\psi$  of  $G - R - v$  extending  $\psi_0$ , we have

$$E[a(G, \varphi, v) | \varphi \text{ extends } \psi] \geq 2\ell.$$

Let  $A$  be the set of colors that  $\psi$  does not use on the neighbors of  $v$ ; since  $v$  is  $\psi$ -rich, we have  $|A| \geq k/2$ . For each  $u \in R$ , let  $L_u$  be the set of colors not used by  $\psi$  on the neighbors of  $u$ . For  $c \in A$ , let  $X_c$  denote the event that  $\varphi$

does not use  $c$  on the neighbors of  $v$ . Observe that

$$\begin{aligned}
& E[a(G, \varphi, v) | \varphi \text{ extends } \psi] \\
&= \sum_{c \in A} \Pr[X_c | \varphi \text{ extends } \psi] \\
&= \sum_{c \in A} \prod_{u \in R: c \in L_u} (1 - 1/|L_u|) \\
&\geq |A| \left( \prod_{c \in A} \prod_{u \in R: c \in L_u} (1 - 1/|L_u|) \right)^{1/|A|} && \text{(AG inequality)} \\
&= |A| \left( \prod_{u \in R} \prod_{c \in L_u \cap A} (1 - 1/|L_u|) \right)^{1/|A|} \\
&\geq |A| \left( \prod_{u \in R} (1 - 1/|L_u|)^{|L(u)|} \right)^{1/|A|} \\
&\geq |A| \left( \prod_{u \in R} (1 - 1/t)^t \right)^{1/|A|} \geq |A| (1 - 1/t)^{t\Delta/|A|} && \text{(since } |L_u| \geq t \text{ and } |R| \leq \Delta) \\
&\geq |A| 4^{-\Delta/|A|} \geq \frac{k}{2} 4^{-2\Delta/k} \geq \frac{k}{2 \cdot 4^{\frac{1}{2} \log \Delta}} && \text{(since } t \geq 2 \text{ and } |A| \geq k/2) \\
&\geq 2 \frac{\Delta^{1-\log 2}}{\log \Delta} = 2\ell.
\end{aligned}$$

This finishes the proof.  $\square$

**Corollary 4.** *If  $G$  is a triangle-free graph with  $n$  vertices and maximum degree  $\Delta$ , then  $G$  has an independent set of size*

$$\Omega \left( \max(n \log \Delta / \Delta, \sqrt{n \log n}) \right).$$

Hence, the Ramsey number  $R(3, m)$  is  $O(\frac{m^2}{\log m})$ .

*Proof.* Since  $G$  has chromatic number  $O(\Delta / \log \Delta)$ , the largest color class has size  $\Omega(n \log \Delta / \Delta)$ . If  $\Delta = \Omega(\sqrt{n \log n})$ , then the neighborhood of a vertex of maximum degree is an independent set of size  $\Omega(\sqrt{n \log n})$ . Otherwise,  $\Delta = O(\sqrt{n \log n})$  and  $\Omega(n \log \Delta / \Delta) = \Omega(\sqrt{n \log n})$ . Hence, if  $n \geq c \frac{m^2}{\log m}$  for a sufficiently large constant  $c$ , then  $G$  has an independent set of size  $\Omega(\sqrt{n \log n}) \geq m$ .  $\square$

All these bounds are tight (but proving this is non-trivial).