

# List coloring and Gallai trees

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## 1 List coloring and choosability

A *list assignment* for a graph  $G$  is a function  $L$  that to each vertex  $v \in V(G)$  assigns a *list*  $L(v)$  of colors. An  $L$ -coloring of  $G$  is a proper coloring  $\varphi$  such that  $\varphi(v) \in L(v)$  for all  $v \in V(G)$ . The *choosability*  $\chi_l(G)$  of  $G$  is the minimum integer  $k$  such that  $G$  can be  $L$ -colored for every assignment  $L$  of lists of size at least  $k$ .

**Observation 1.** *Every graph  $G$  satisfies*

$$\chi(G) \leq \chi_l(G).$$

*If  $G$  is  $d$ -degenerate, then*

$$\chi_l(G) \leq d + 1.$$

Choosability matches the chromatic number for many graphs (cycles, cliques, ...). However, the choosability can be arbitrarily large compared to the chromatic number, as the following result shows.

**Lemma 2.** *For every positive integer  $a$ ,*

$$\chi_l(K_{a,a^a}) = a + 1.$$

*Proof.* For every  $n$ , the bipartite graph  $K_{a,n}$  is  $a$ -degenerate, implying that  $\chi_l(K_{a,n}) \leq a + 1$ . Hence, it suffices to show that there exists an assignment of lists of size  $a$  to  $K_{a,a^a}$  from that the graph cannot be colored.

Let  $A$  and  $B$  be the parts of  $K_{a,a^a}$ , where  $|A| = a$  and  $|B| = a^a$ . Let  $A = \{v_1, \dots, v_a\}$ . Since  $|B|$  is equal to the number of sequences of numbers  $\{1, \dots, a\}$  of length  $a$ , we can label vertices of  $B$  as  $u_{i_1, \dots, i_a}$  for  $1 \leq i_1, \dots, i_a \leq a$ . Let us give vertices of  $A$  pairwise disjoint lists, say  $L(v_i) = \{(i, 1), (i, 2), \dots, (i, a)\}$  for  $1 \leq i \leq a$ . We give vertices of  $B$  different lists, each of them intersecting the list of each vertex of  $A$  in exactly one

color; say  $L(u_{i_1, \dots, i_a}) = \{(1, i_1), (2, i_2), \dots, (a, i_a)\}$  for  $1 \leq i_1, \dots, i_a \leq a$ . We claim that  $K_{a, a^a}$  is not  $L$ -colorable; indeed, if we give vertices of  $A$  colors  $(1, c_1), (2, c_2), \dots, (a, c_a)$ , then all colors in the list of the vertex  $u_{c_1, c_2, \dots, c_a}$  are used on its neighbors, and thus this vertex cannot be colored.  $\square$

On the other hand, we have the following positive result on choosability of bipartite graphs.

**Lemma 3.** *For every positive integer  $n$ ,*

$$\chi_l(K_{n,n}) \leq \lfloor \log_2 n \rfloor + 2.$$

*Proof.* Let  $c = \lfloor \log_2 n \rfloor + 2$  and let  $L$  be any assignment of lists of size at least  $c$  to vertices of  $K_{n,n}$ . Let  $A$  and  $B$  be the parts of  $K_{n,n}$ . For each color, we flip a fair coin and according to the result we delete it either from the lists of all vertices of  $A$  or from the lists of all vertices of  $B$ . Afterwards, the lists of vertices of  $A$  are disjoint from the lists of vertices of  $B$ , and thus if these lists are non-empty, we can properly color  $K_{n,n}$ . The probability that a list of a vertex  $v \in V(K_{n,n})$  becomes empty is at most  $2^{-c}$ . Hence, the expected number of empty lists is at most

$$2^{-c}|V(K_{n,n})| = \frac{2n}{2^{\lfloor \log_2 n \rfloor + 2}} < 1.$$

Hence, with non-zero probability, all the lists are non-empty, and thus  $K_{n,n}$  can be  $L$ -colored. Since the choice of  $L$  was arbitrary, it follows that  $\chi_l(K_{n,n}) \leq c$  as required.  $\square$

Note that

$$\chi_l(K_{n,n}) \geq \frac{\log_2 n}{\log_2 \log_2 n}$$

by Lemma 2; a more involved argument shows that actually  $\chi_l(K_{n,n}) = \Omega(\log n)$ .

## 2 Planar graphs

By Observation 1, all planar graphs are 6-choosable. Thomassen proved that they are actually 5-choosable, strengthening the 5-color theorem. In fact, he proved the following stronger claim.

**Theorem 4.** *Let  $G$  be a plane graph, let  $p_1 p_2$  be an edge contained in the boundary of its outer face, and let  $L$  be a list assignment for  $G$  satisfying the following.*

- $|L(v)| \geq 5$  for every vertex  $v \in V(G)$  not incident with the outer face.
- $|L(v)| \geq 3$  for every vertex  $v \in V(G)$  incident with the outer face and distinct from  $p_1$  and  $p_2$ .
- $|L(p_1)|, |L(p_2)| \geq 1$ , and if  $|L(p_1)| = |L(p_2)| = 1$ , then  $L(p_1) \neq L(p_2)$ .

Then  $G$  is  $L$ -colorable.

*Proof.* We proceed by induction on the number of vertices of  $G$ . The case  $|V(G)| = 2$  is trivial, hence assume  $|V(G)| \geq 3$ .

We can assume  $G$  is connected, otherwise we apply induction to each component of  $G$ . Furthermore, we can assume  $G$  is 2-connected. Otherwise,  $G = G_1 \cup G_2$  for proper induced subgraphs  $G_1$  and  $G_2$  of  $G$  intersecting in exactly one vertex  $v$ . We can assume that  $p_1 p_2 \in E(G_1)$ . By the induction hypothesis, there exists an  $L$ -coloring  $\varphi_1$  of  $G_1$ . Let  $L'(v) = \{\varphi_1(v)\}$  and  $L'(x) = L(x)$  for all  $x \in V(G_2) \setminus \{v\}$ . Then  $G_2$  with the list assignment  $L'$  satisfies the assumptions of the theorem (with  $v$  and one of its neighbors playing the role of  $p_1 p_2$ ), and thus  $G_2$  has an  $L'$ -coloring  $\varphi_2$ . The colorings  $\varphi_1$  and  $\varphi_2$  together give an  $L$ -coloring of  $G$ .

Since  $G$  is 2-connected, its outer face is bounded by a cycle  $C$ . We can assume that the cycle  $C$  is induced. Otherwise, if  $C$  has a chord  $v_1 v_2$ , then  $G = G_1 \cup G_2$  for proper induced subgraphs  $G_1$  and  $G_2$  of  $G$  intersecting exactly in  $v_1 v_2$ . We can assume that  $p_1 p_2 \in E(G_1)$ . By the induction hypothesis, there exists an  $L$ -coloring  $\varphi_1$  of  $G_1$ . Let  $L'(v_1) = \{\varphi_1(v_1)\}$ ,  $L'(v_2) = \{\varphi_1(v_2)\}$  and  $L'(x) = L(x)$  for all  $x \in V(G_2) \setminus \{v_1, v_2\}$ . Then  $G_2$  with the list assignment  $L'$  satisfies the assumptions of the theorem (with  $v_1 v_2$  playing the role of  $p_1 p_2$ ), and thus  $G_2$  has an  $L'$ -coloring  $\varphi_2$ . The colorings  $\varphi_1$  and  $\varphi_2$  together give an  $L$ -coloring of  $G$ .

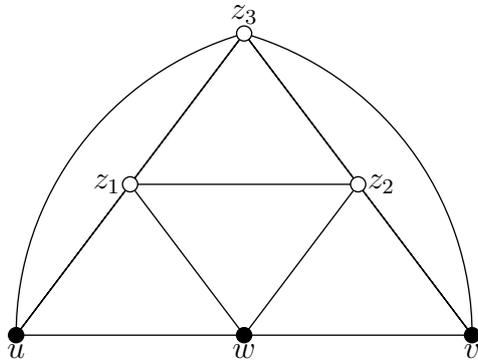
We can also assume that  $|L(p_1)| = |L(p_2)| = 1$ , as otherwise we can throw away extra colors from their lists. Let  $C = p_1 p_2 p_3 \dots p_k$ . By the assumptions,  $|L(p_3)| \geq 3$ , and thus there exist two distinct colors  $c_1, c_2 \in L(p_3) \setminus L(p_2)$ . Let  $t \geq 3$  be the maximum integer such that  $\{c_1, c_2\} \subseteq L(p_i)$  for  $3 \leq i \leq t$ . Let  $v = p_{t+1}$  if  $t < k$  and  $v = p_1$  if  $t = k$ . By the choice of  $t$  and since  $|L(p_1)| = 1$ , we can assume that  $c_2 \notin L(v)$ . Let  $G' = G - \{p_3, \dots, p_t\}$  and let  $L'$  be the list assignment for  $G'$  such that  $L'(p_2) = L(p_2)$ ,  $L'(v) = L(v)$ ,  $L'(x) = L(x)$  if  $x \in V(G') \setminus \{v, p_2\}$  has no neighbor in  $\{p_3, \dots, p_t\}$ , and  $L'(x) = L(x) \setminus \{c_1, c_2\}$  otherwise. Since the outer face of  $G$  is bounded by an induced cycle  $C$ , if  $x \in V(G') \setminus \{v, p_2\}$  has a neighbor in  $\{p_3, \dots, p_t\}$ , then  $x \notin V(C)$ , and thus  $|L(x)| \geq 5$  and  $|L'(x)| \geq 3$ ; furthermore, such a vertex  $x$  is contained in the boundary of the outer face of  $G'$ . We conclude that  $G'$  with the list assignment  $L'$  satisfies the assumptions of the theorem,

and thus  $G'$  has an  $L'$ -coloring  $\varphi'$  by the induction hypothesis. By the choice of  $L'$ , none of the neighbors of  $\{p_3, \dots, p_t\}$  in  $G'$  except for  $v$  can be given color  $c_1$  or  $c_2$ . Recall also that  $c_2 \notin L(v)$ , and thus  $\varphi'(v) \neq c_2$ . Hence, we can extend  $\varphi'$  to an  $L$ -coloring of  $G$  by giving  $p_t, p_{t-2}, \dots$  the color  $c_2$  and  $p_{t-1}, p_{t-3}, \dots$  the color  $c_1$ .  $\square$

However, in contrast to the Four Color Theorem, not all planar graphs are 4-choosable.

**Lemma 5.** *There exists a planar graph  $G$  that is not 4-choosable.*

*Proof.* Let  $G_{uvw}$  be the following graph.



Let  $L_{a,m,b}$  (with distinct  $a, m, b \notin \{11, 12\}$ ) be the list assignment such that  $L_{a,m,b}(z_1) = \{a, m, 11, 12\}$ ,  $L_{a,m,b}(z_2) = \{m, b, 11, 12\}$ , and  $L_{a,m,b}(z_3) = \{a, b, 11, 12\}$ . Then a precoloring of  $(u, w, v)$  by colors  $(a, m, b)$  cannot be extended to an  $L_{a,m,b}$ -coloring of  $G_{uvw}$ .

Let  $G_{uv}$  be the graph formed by two copies of  $G_{uvw}$  sharing the path  $uvw$ . Let  $L_{a,b}$  (with distinct  $a, b \notin \{9, 10, 11, 12\}$ ) be the list assignment matching  $L_{a,9,b}$  in one of the copies,  $L_{a,10,b}$  in the other copy, and with  $L_{a,b}(w) = \{a, b, 9, 10\}$ . Then a precoloring of  $(u, v)$  by colors  $\{a, b\}$  cannot be extended to an  $L_{a,b}$ -coloring of  $G_{uv}$ .

Let  $G$  be the graph formed by 16 copies of  $G_{uv}$  sharing the vertices  $u$  and  $v$ . Let  $L(u) = \{1, 2, 3, 4\}$ ,  $L(v) = \{5, 6, 7, 8\}$ , and let  $L$  match  $L_{a,b}$  for  $a \in \{1, 2, 3, 4\}$  and  $b \in \{5, 6, 7, 8\}$  on the 16 copies of  $G_{uv}$ . Then  $G$  is not  $L$ -colorable.  $\square$

### 3 Degree choosability

We want to obtain a list version of Brooks' theorem.

**Theorem 6** (Brooks). *Let  $G$  be a connected graph of maximum degree at most  $\Delta$ . If  $G$  is not  $\Delta$ -colorable, then either  $G = K_{\Delta+1}$ , or  $\Delta = 2$  and  $G$  is an odd cycle.*

A *degree assignment* to a graph  $G$  is a list assignment such that  $|L(v)| \geq \deg(v)$  for all  $v \in V(G)$ .

**Lemma 7.** *Let  $G$  be a connected graph and let  $L$  be a degree assignment for  $G$ . If  $G$  is not  $L$ -colorable, then  $|L(v)| = \deg(v)$  for all  $v \in V(G)$ .*

*Proof.* If  $|L(v)| > \deg(v)$ , then let  $v_1, \dots, v_n$  be a listing of vertices of  $G$  in non-increasing order according to their distance from  $v$ ; hence,  $v_n = v$  and for  $1 \leq i \leq n-1$ , the vertex  $v_i$  has a neighbor  $v_j$  with  $j > i$  (the neighbor of  $v_i$  on a shortest path from  $v_i$  to  $v$ ). Let us greedily  $L$ -color  $v_1, \dots, v_n$  in order. For  $1 \leq i \leq n-1$ , at least one neighbor of  $v_i$  has not been colored yet, and thus at most  $\deg(v_i) - 1 < |L(v_i)|$  colors need to be avoided. At  $v_n$ , at most  $\deg(v_n) < |L(v_n)|$  colors need to be avoided. Hence, in both cases, we can give  $v_i$  a color from its list different from the colors of its neighbors.  $\square$

**Corollary 8.** *Let  $G$  be a connected graph and let  $L$  be a degree assignment for  $G$ . If  $G$  is not  $L$ -colorable,  $uv \in E(G)$  and  $u$  is not a cutvertex in  $G$ , then  $L(u) \subseteq L(v)$ .*

*Proof.* Otherwise, there exists a color  $c \in L(u) \setminus L(v)$ . Let  $G' = G - u$ ; since  $u$  is not a cutvertex,  $G'$  is connected. Let  $L'(x) = L(x) \setminus \{c\}$  for all neighbors  $x$  of  $u$  and  $L'(x) = L(x)$  for all non-neighbors  $x$ . Note that the list size decreases only for neighbors  $x$  of  $u$  for which  $\deg_{G'}(x) = \deg_G(x) - 1$ , and thus  $L'$  is a degree assignment for  $G'$ . Furthermore,  $|L'(v)| = |L(v)| \geq \deg_G(v) > \deg_{G'}(v)$ , and thus  $G'$  is  $L'$ -colorable by Lemma 7. We can extend this coloring to an  $L$ -coloring of  $G$  by giving  $v$  color  $c$ .  $\square$

Note that if neither  $u$  nor  $v$  is a cutvertex, then Corollary 8 implies  $L(u) \subseteq L(v)$  and  $L(v) \subseteq L(u)$ , and thus  $L(u) = L(v)$ .

**Corollary 9.** *Let  $G$  be a 2-connected graph and let  $L$  be a degree assignment for  $G$ . Then  $G$  is not  $L$ -colorable if and only if  $G$  is a clique or an odd cycle and all vertices of  $G$  have the same list of length equal to the degree of vertices of  $G$ .*

*Proof.* The “if” part is trivial. For the “only if” part, suppose that  $G$  is not  $L$ -colorable. Since  $G$  is 2-connected, Corollary 8 implies that any two adjacent vertices of  $G$  have the same list, and consequently all the vertices of  $G$  have the same list, say  $\{1, \dots, d\}$ , where  $d \leq \Delta(G)$ . It follows that  $G$  is not  $d$ -colorable, and thus either  $G = K_{d+1}$  or  $d = 2$  and  $G$  is an odd cycle by Theorem 6.  $\square$

A *Gallai tree* is a connected graph  $T$  such that every 2-connected block of  $T$  is either a clique or an odd cycle. Suppose  $B_1, \dots, B_k$  are the blocks of a Gallai tree  $T$ , and let  $S_1, \dots, S_k$  be sets of colors satisfying the following conditions:

- For  $1 \leq i \leq k$ , if  $B_i$  is a clique, then  $|S_i| = |V(B_i)| - 1$ , and if  $B_i$  is an odd cycle, then  $|S_i| = 2$ .
- For  $1 \leq i < j \leq k$ , if  $B_i \cap B_j \neq \emptyset$ , then  $S_i \cap S_j = \emptyset$ .

For  $v \in V(T)$ , let  $L(v) = \bigcup_{v \in B_i} S_i$ . If a list assignment  $L$  can be expressed in this way, we say that  $L$  is a *blockwise uniform* assignment for  $T$ .

**Theorem 10** (Gallai). *Let  $G$  be a connected graph and let  $L$  be a degree assignment for  $G$ . Then  $G$  is not  $L$ -colorable if and only if  $G$  is a Gallai tree and  $L$  is blockwise uniform.*

*Proof.* It is easy to see that a Gallai tree cannot be colored from a blockwise uniform assignment, and thus it suffices to prove the “only if” part. We do the proof by induction on the number of vertices of  $G$ .

By Corollary 9, the claim holds when  $G$  is 2-connected. Hence, suppose that  $G$  is not 2-connected. First, we prove that  $G$  is a Gallai tree. Let  $B$  be a block of  $G$ . Since  $G$  is not 2-connected, there exists a leaf block  $B'$  of  $G$  distinct from  $B$ . Let  $v$  be a vertex of  $B'$  which is not a cutvertex, and let  $G' = G - v$ . Let  $c$  be any color in  $L(v)$  and let  $L'(x) = L(x) \setminus \{c\}$  for all neighbors  $x$  of  $v$  and  $L'(x) = L(x)$  for all other vertices  $x$  of  $G'$ . Note that  $L'$  is a degree assignment for  $G'$  and that  $G'$  is not  $L'$ -colorable, as otherwise we can extend the coloring to an  $L$ -coloring of  $G$  by giving  $v$  the color  $c$ . By the induction hypothesis,  $G'$  is a Gallai tree. Note that  $B$  is also a block of  $G'$ , and thus  $B$  is a clique or an odd cycle. As the choice of  $B$  was arbitrary, all blocks of  $G$  are cliques or odd cycles, and thus  $G$  is a Gallai tree.

Let  $B_1, \dots, B_k$  be the blocks of  $G$ , where without loss of generality  $B_k$  is a leaf block. Let  $z$  be the cutvertex of  $B_k$  and let  $v$  be any other vertex of  $B_k$ , and let  $S_k = L(v)$ ; by Corollary 8, we conclude that all non-cut vertices of  $B_k$  have list  $S_k$ , and  $S_k \subseteq L(z)$ . By Lemma 7, if  $B_k$  is a clique then  $|S_k| = |B_k| - 1$ , and if  $B_k$  is an odd cycle, then  $|S_k| = 2$ . Let  $G' = B_1 \cup \dots \cup B_{k-1}$ , let  $L'(x) = L(x)$  for  $x \in V(G') \setminus \{z\}$  and  $L'(z) = L(z) \setminus S_k$ . Note that  $L'$  is a degree assignment for  $G'$  and that  $G'$  is not  $L'$ -colorable, as otherwise the coloring would extend to an  $L$ -coloring of  $G$  by using the colors in  $S_k$  to color  $B_k - z$ . By the induction hypothesis,  $L'$  is blockwise uniform as shown by sets  $S_1, \dots, S_{k-1}$ . But then the sets  $S_1, \dots, S_{k-1}, S_k$  show that  $L$  is blockwise uniform.  $\square$

**Corollary 11.** *Let  $G$  be a  $(c+1)$ -critical graph and let  $S$  be the set of vertices of  $G$  of degree  $c$ . Then each component of  $G[S]$  is a Gallai tree.*

*Proof.* Consider any component  $C$  of  $G[S]$ . Since  $G$  is  $(c+1)$ -critical,  $G - C$  has a  $c$ -coloring  $\varphi$ . Let  $L$  be the list assignment to  $G[C]$  in which for each  $v \in C$ , the list  $L(v)$  consists of those of colors  $\{1, \dots, c\}$  that are not used by  $\varphi$  on the neighbors of  $v$ . If  $v$  has  $k$  neighbors in  $V(G) \setminus V(C)$ , then  $\deg_{G[C]}(v) = \deg_G(v) - k = c - k$  and  $|L(v)| \geq c - k$ , and thus  $L$  is a degree assignment for  $G[C]$ . An  $L$ -coloring of  $G[C]$  together with  $\varphi$  would give a  $c$ -coloring of  $G$ ; since  $G$  is  $(c+1)$ -critical, we conclude that  $G[C]$  is not  $L$ -colorable, and by Theorem 10,  $G[C]$  is a Gallai tree.  $\square$