

200 mil. Kč

200

mil / student

20 years

50k / student / year

offline lecture -

- speed-up video (150%)
- max (30%) extension
- pause, skip backward

* exams - 1/2 half of Term.

Communication complexity

$$EQ(x, y) = [x \stackrel{?}{=} y]$$

$$\text{protocol} \leq \underline{\underline{n+1}}$$

\vee
 $n+1$

• randomized protocol?

- pick a random bit i and compare
 Alice $i, x_i \rightarrow$ Bob
 $\leftarrow y_i$

$x=y$
 $x \neq y$

Pr[error] = 0

$$Pr[\text{error}] = 1 - \frac{O_{Ham}(x, y)}{n} = \underline{\underline{1 - \frac{1}{n}}}$$

repeat l times

$$Pr[\text{error}] \leq \left(1 - \frac{1}{n}\right)^l \leq e^{-\frac{l}{n}}$$

$$l = n$$

$$\left(1 - \frac{1}{n}\right)^n \approx \frac{1}{e}$$

$$(1-x)^l \leq e^{-xl}$$

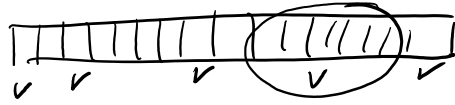
$$\underline{\underline{l \approx n}}$$

$l \approx n$

• computing parities for random subset of bits

Alice \rightarrow Bob

$\sum x_i \pmod 2$



$x = y$
 $x \neq y$

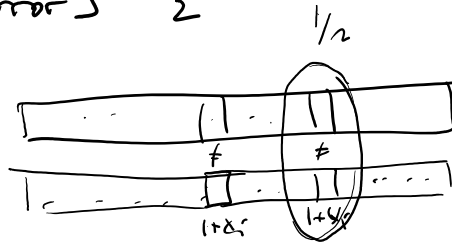
$\Pr[\text{error}] = 0$

$\Pr[\text{error}] = \frac{1}{2}$

$x_i = 1 \oplus y_i$

$x_i \neq y_i \Leftrightarrow$

$\sum y_i = \sum 1 \oplus x_i$
 $= \sum 1 + \sum x_i$



repeat l -times

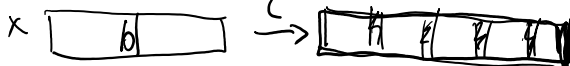
$\Pr[\text{error}] = \left(\frac{1}{2}\right)^l = \frac{1}{2^l}$
 $= \frac{1}{n^2}$

$l=10$ 99.9%

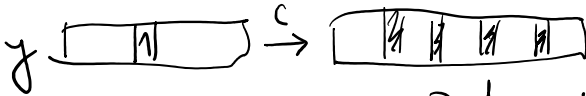
(shared random string)

$a = 00010000$

$C [N, n, d]$



$C(x) = x' \quad N$
 $C(y) = y' \quad N \geq \frac{d}{N}$



$\Pr[\text{error}] \leq 1 - \frac{d}{N}$

$\left[\frac{(1-\epsilon)}{2}n, n, \epsilon n \right] \quad \delta = H(\epsilon)$

$\Pr[\text{error}] \leq 1 - \epsilon$

$\lg N = \lg((1-\epsilon)n) = \underline{\underline{O(\lg n)}}$

repeat l -times

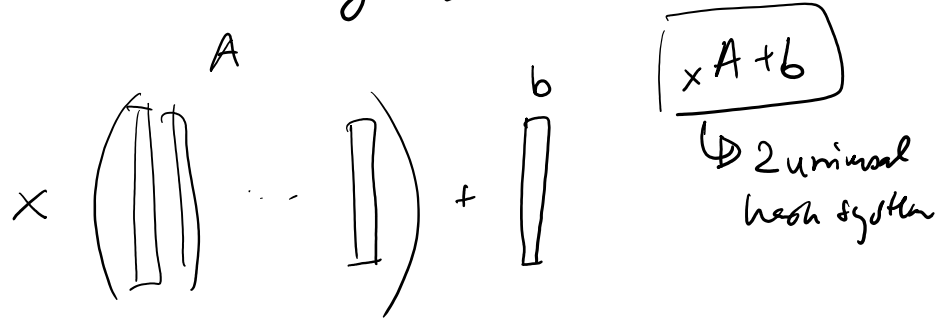
$\Pr[\text{error}] \leq (1-\epsilon)^l$

[doesn't use shared random bits]

$O(\lg n) \quad O(\lg n) \quad \frac{1}{n^2}$

→ $O(\lg n)$ bits

- hashing \approx error correcting codes

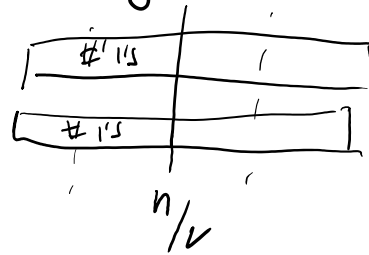


• $MED(x, y) = \text{median of } x \cup y$

$y, x \subseteq \{1, \dots, n\}$

y has odd #'s
x has odd even #'s

- binary search



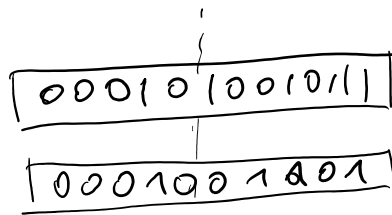
$O(\lg^2 n)$ in total

• $GT(x, y) = [x > y]$

0001
1 0101

n+1

2 GT \geq EQ



• binary search on to find the differing position using EQ protocol. A

Q: • deterministic? $\geq \frac{n}{2}$

$\lg n$
runs

EQ protocol using $O(\lg^2 n)$ bits

error prob $\leq \frac{1}{n^2}$

→ error prob of the GT-protocol $\leq \frac{\lg n}{n^2}$

err prob of single check $\leq \frac{1}{2^n}$

$$\text{error prob.} \leq \frac{1}{n^2}$$
$$\rightarrow \text{error prob. of the GT-protocol} \leq \frac{\lg n}{n^2}$$

$$O(\lg^3 n) \text{ bits.}$$

$$\text{check} \leq \frac{1}{\lg n}$$

$$\left\{ O(\lg n) \text{ bits} \right\}$$