## Proof of Regularity Lemma

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We are going to use the following well-known inequality.

**Lemma 1** (Cauchy-Schwarz inequality). For any real numbers  $u_1, \ldots, u_n$ and  $v_1, \ldots, v_n$ , we have  $\left(\sum_{i=1}^n u_i^2\right) \left(\sum_{i=1}^n v_i^2\right) \ge \left(\sum_{i=1}^n u_i v_i\right)^2$ .

Our goal is to prove the Regularity Lemma.

**Theorem 2.** For any positive integer  $m_0$  and real number  $\varepsilon > 0$ , there exists an integer  $M \ge m_0$  such that the following holds. Every graph G with at least  $m_0$  vertices has an  $\varepsilon$ -regular partition of order at least  $m_0$  and at most M.

Without loss of generality, we can assume that  $\varepsilon < 1/2$ . Let  $n \ge m_0$  be the number of vertices of G. Let us recall that an  $\varepsilon$ -regular partition consists of parts  $V_1, \ldots, V_m$  of the same size, and an exceptional part  $V_0$  of size at most  $\varepsilon n$ . For the purposes of the proof, it will be more convenient to further break up  $V_0$  to single-element parts. Hence, we say that a partition  $\mathcal{P}$  of V(G) is an  $\varepsilon$ -regular partition of order m if  $\mathcal{P}$  contains distinct parts  $V_1, \ldots, V_m$  such that

- (i)  $|V_1| = |V_2| = \ldots = |V_m|,$
- (ii) for all but at most  $\varepsilon m^2$  values of  $1 \leq i < j \leq m$ , the pair  $(V_i, V_j)$  is  $\varepsilon$ -regular, and
- (iii)  $\mathcal{P} \setminus \{V_1, \ldots, V_m\}$  consists of at most  $\varepsilon n$  parts of size one.

The proof proceeds by starting with an initial partition of order  $m_0$  satisfying (i) and (iii) and proceeds to refine it, increasing a suitably defined "quality" of the partition, so that the following conditions hold:

• each refinement increases the quality by at least  $\frac{\varepsilon^5}{4}n^2$ ,

- each refinement increases the order of the partition at most exponentially, and
- once further refinement is not possible, (ii) holds.

Moreover, the quality is bounded from above by  $n^2$ , guaranteeing that the process ends after at most  $\frac{4}{\varepsilon^5}$  iterations, at which point the order is bounded by the tower function of exponentials of height  $\frac{4}{\varepsilon^5}$  and ending in  $m_0$ , i.e., a constant M depending only on  $\varepsilon$  and  $m_0$ .

Let us now show the details of this argument, starting with the definition of the quality function. For disjoint sets  $A, B \subseteq V(G)$ , let  $q(A, B) = |A||B|d^2(A, B)$ . For a partition  $\mathcal{P} = \{P_1, \ldots, P_k\}$  of V(G), let

$$q(\mathcal{P}) = \sum_{i < j} q(P_i, P_j).$$

Note that  $q(A, B) \leq |A||B|$  and  $q(\mathcal{P}) \leq \sum_{i < j} |P_i||P_j| < n^2$ . Next, let us show that the quality does not decrease when we refine the partition.

**Lemma 3.** If  $\{A_1, \ldots, A_s\}$  is a partition of A and  $\{B_1, \ldots, B_s\}$  is a partition of B, then  $\sum_{i,j} q(A_i, B_j) \ge q(A, B)$ .

*Proof.* We have  $e(A, B) = \sum_{i,j} e(A_i, B_j)$  and  $|A||B| = \sum_{i,j} |A_i||B_j|$ . Let  $u_{i,j} = \sqrt{|A_i||B_j|}d(A_i, B_j)$  and  $v_{i,j} = \sqrt{|A_i||B_j|}$ . Therefore,

$$\sum_{i,j} u_{i,j}^2 = \sum_{i,j} q(A_i, B_j)$$
  
$$\sum_{i,j} v_{i,j}^2 = \sum_{i,j} |A_i| |B_j| = |A| |B|$$
  
$$\sum_{i,j} u_{i,j} v_{i,j} = \sum_{i,j} d(A_i, B_j) |A_i| |B_j| = \sum_{i,j} e(A_i, B_j) = e(A, B).$$

Cauchy-Schwarz inequality gives

$$|A||B|\sum_{i,j}q(A_i, B_j) \ge e^2(A, B) = |A|^2|B|^2d^2(A, B),$$

and thus

$$\sum_{i,j} q(A_i, B_j) \ge |A| |B| d^2(A, B) = q(A, B).$$

We say that a partition  $\mathcal{P}'$  is a *refinement* of a partition  $\mathcal{P}$  if every part of  $\mathcal{P}'$  is a subset of a part of  $\mathcal{P}$ . The *common refinement* of two partitions  $\mathcal{P}_1$  and  $\mathcal{P}_2$  is the partition  $\{A \cap B : A \in \mathcal{P}_1, B \in \mathcal{P}_2, A \cap B \neq \emptyset\}$ .

**Corollary 4.** If a partition  $\mathcal{P}'$  is a refinement of a partition  $\mathcal{P}$ , then  $q(\mathcal{P}') \geq q(\mathcal{P})$ .

*Proof.* Let  $\mathcal{P} = \{P_1, \ldots, P_s\}$ . For  $1 \leq i \leq s$ , let  $\mathcal{P}_i = \{Q \in \mathcal{P}' : Q \subseteq P_i\}$ , so that  $\mathcal{P}' = \mathcal{P}_1 \cup \ldots \cup \mathcal{P}_s$ . Then

$$q(\mathcal{P}') \geq \sum_{1 \leq i < j \leq s} \sum_{Q \in \mathcal{P}_i, R \in \mathcal{P}_j} q(Q, R) \geq \sum_{1 \leq i < j \leq s} q(P_i, P_j) = q(\mathcal{P})$$

by Lemma 3.

We also need to know that if (A, B) is not an  $\varepsilon$ -pair, then we can partition it with a substantial increase in quality.

**Lemma 5.** Suppose (A, B) is a non- $\varepsilon$ -regular pair of disjoint subsets of vertices of a graph G; i.e., there exist sets  $A_1 \subseteq A$  and  $B_1 \subseteq B$  such that  $|A_1| \ge \varepsilon |A|, |B_1| \ge \varepsilon |B|, \text{ and } |d(A_1, B_1) - d(A, B)| > \varepsilon$ . For  $A_2 = A \setminus A_1$  and  $B_2 = B \setminus B_1$ , we have  $\sum_{1 \le i, j \le 2} q(A_i, B_j) > q(A, B) + \varepsilon^4 |A| |B|$ .

*Proof.* Let d = d(A, B). For  $1 \le i, j \le 2$ ,  $u \in A_i$ , and  $v \in B_j$ , let  $d(u, v) = d(A_i, B_j)$  and  $\delta(u, v) = d(u, v) - d$ . Note that  $|\delta(u, v)| > \varepsilon$  for  $u \in A_1$  and  $v \in B_1$ . We have

$$\begin{aligned} d|A||B| &= e(A,B) = \sum_{1 \le i,j \le 2} e(A_i,B_j) \\ &= \sum_{1 \le i,j \le 2} d(A_i,B_j)|A_i||B_j| = \sum_{u \in A, v \in B} d(u,v) \\ &= d|A||B| + \sum_{u \in A, v \in B} \delta(u,v), \end{aligned}$$

and thus

$$\sum_{u \in A, v \in B} \delta(u, v) = 0.$$

Moreover,  $q(A, B) = |A||B|d^2$  and

$$\sum_{1 \le i,j \le 2} q(A_i, B_j) = \sum_{1 \le i,j \le 2} |A_i| |B_j| d^2(A_i, B_j) = \sum_{u \in A, v \in B} d^2(u, v)$$
  
= 
$$\sum_{u \in A, v \in B} (d + \delta(u, v))^2 = \sum_{u \in A, v \in B} (d^2 + 2d\delta(u, v) + \delta^2(u, v))$$
  
= 
$$|A| |B| d^2 + 2d \sum_{u \in A, v \in B} \delta(u, v) + \sum_{u \in A, v \in B} \delta^2(u, v)$$
  
= 
$$q(A, B) + \sum_{u \in A, v \in B} \delta^2(u, v) \ge q(A, B) + \sum_{u \in A_1, v \in B_1} \delta^2(u, v)$$
  
> 
$$q(A, B) + |A_1| |B_1| \varepsilon^2 \ge q(A, B) + \varepsilon^4 |A| |B|.$$

Let  $\mathcal{P} = \{V_1, \ldots, V_s, V_{s+1}, \ldots, V_{s+t}\}$  be a partition of V(G) such that  $t \leq \varepsilon n$ ,  $|V_i| = 1$  for  $i \geq s+1$  and  $|V_1| = \ldots = |V_s| = N$ . Note that  $sN = n - t \geq (1 - \varepsilon)n \geq n/2$ . Suppose that  $\mathcal{P}$  is not an  $\varepsilon$ -regular partition of order s, i.e., that it does not satisfy (ii). Let X be the set of pairs (i, j) such that  $1 \leq i < j \leq s$  and  $(V_i, V_j)$  is not  $\varepsilon$ -regular; we have  $|X| \geq \varepsilon s^2$ . For each  $(i, j) \in X$ , let  $\mathcal{P}_{ij} = \{V_{i,j,1}, V_{i,j,2}, V_{j,i,1}, V_{j,i,2}, V(G) \setminus (V_i \cup V_j)\}$  be the partition of V(G) where

$$V_{i} = V_{i,j,1} \cup V_{i,j,2}, V_{j} = V_{j,i,1} \cup V_{j,i,2},$$
$$|V_{i,j,1}| \ge \varepsilon |V_{i}|, |V_{j,i,1}| \ge \varepsilon |V_{j}|, \text{ and }$$
$$d(V_{i,j,1}, V_{j,i,1}) - d(V_{i}, V_{j})| > \varepsilon.$$

Let  $\mathcal{P}'$  be the common refinement of  $\mathcal{P}$  and all partitions  $\mathcal{P}_{ij}$  for  $(i, j) \in X$ . Let  $\mathcal{V}_i = \{A \in \mathcal{P}' : A \subseteq V_i\}$ . We have

$$q(\mathcal{P}') \ge \sum_{i < j} \sum_{A \in \mathcal{V}_i, B \in \mathcal{V}_j} q(A, B).$$

Lemma 3 implies

$$\sum_{A \in \mathcal{V}_i, B \in \mathcal{V}_j} q(A, B) \ge q(V_i, V_j)$$

for every i < j. Moreover, if  $(i, j) \in X$ , then

$$\sum_{A \in \mathcal{V}_i, B \in \mathcal{V}_j} q(A, B) \ge \sum_{1 \le o, p \le 2} q(V_{i,j,o}, V_{j,i,p}) > q(V_i, V_j) + \varepsilon^4 |V_i| |V_j|$$

by Lemma 5. We conclude that

$$q(\mathcal{P}') > \left(\sum_{i < j} q(V_i, V_j)\right) + |X|\varepsilon^4 N^2 \ge q(\mathcal{P}) + \varepsilon^5 s^2 N^2 \ge q(\mathcal{P}) + \frac{\varepsilon^5}{4}n^2.$$

The partition  $\mathcal{P}'$  has at most  $s2^s$  non-singleton parts (each of the parts  $V_1, \ldots, V_s$  is contained in less than s pairs of X, and thus it is divided less than s times). However, the non-singleton parts can have different sizes. Let  $f(s) = s4^s$ . We divide every non-singleton part of  $\mathcal{P}'$  into as many parts of size  $\left\lceil \frac{N}{4^s} \right\rceil$  as possible, and their remainders to singleton parts. This gives a refinement  $\mathcal{P}''$  of  $\mathcal{P}'$  such that

- every non-singleton part of  $\mathcal{P}''$  has size exactly  $\left\lceil \frac{N}{4^s} \right\rceil$ ,
- the number of non-singleton parts of  $\mathcal{P}''$  is at most f(s), since  $\mathcal{P}$  has s non-singleton parts of size N and at most  $4^s$  parts of size  $\left\lceil \frac{N}{4^s} \right\rceil$  fit into each of them, and
- the number of singleton parts of  $\mathcal{P}''$  is by at most  $s2^s \frac{N}{4^s} \leq s\frac{n/s}{2^s} = \frac{n}{2^s}$  larger than the number of singleton parts of  $\mathcal{P}$ .

Moreover,  $q(\mathcal{P}'') \ge q(\mathcal{P}') > q(\mathcal{P}) + \frac{\varepsilon^5}{4}n^2$ .

Proof of Regularity Lemma. Let

$$s_0 = \max\left(2m_0, \left\lceil \log \frac{8}{\varepsilon^6} \right\rceil\right)$$
, and  
 $M = \max\left(\left\lceil \frac{2}{\varepsilon}s_0 \right\rceil, f(f(\dots f(s_0)\dots))\right)$ ,

where f is iterated  $\left|\frac{4}{\epsilon^5}\right|$ -times.

If G has at most M vertices, then G has an  $\varepsilon$ -regular partition to  $|V(G)| \leq M$  singleton parts. Otherwise, let us start with an arbitrary partition of V(G) to  $s_0$  parts of the same size and less than  $s_0 \leq \frac{\varepsilon}{2}n$  singleton parts. We repeat the refinement procedure described until we reach an  $\varepsilon$ -regular partition  $\mathcal{P}$ . Since  $0 \leq q(\mathcal{Q}) \leq n^2$  for every partition  $\mathcal{Q}$  of V(G) and each iteration of the procedure increases the value of q by more than  $\frac{\varepsilon^5}{4}n^2$ , the number of iterations is at most  $\frac{4}{\varepsilon^5}$ , and thus  $\mathcal{P}$  has at most M non-singleton parts.

We need to verify that the number of singleton parts does not increase over  $\varepsilon n$ . Each iteration increases the number of singleton parts by at most  $\frac{n}{2^s} \leq \frac{n}{2^{s_0}}$ . Therefore, after  $\lfloor \frac{4}{\varepsilon^5} \rfloor$  iterations, there are at most  $\frac{4}{\varepsilon^5 2^{s_0}}n + \frac{\varepsilon}{2}n \leq \varepsilon n$ singleton parts.

Since the initial partition consists of parts of size at most  $n/s_0$ , this is also an upper bound on the size of the parts of one at most  $n/s_0$ , this is of  $\mathcal{P}$  is at least  $\frac{(1-\varepsilon)n}{n/s_0} \ge s_0/2 \ge m_0$ .