Graph polynomials from simple graph sequences

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Graph polynomials Graph homomorphisms

Chromatic polynomial

Definition by evaluations at positive integers

 $k \in \mathbb{N}, \quad P(G; k) = #\{ \text{proper vertex } k \text{-colourings of } G \}.$

$$P(G;k) = \sum_{1 \le j \le |V(G)|} a_j(G) k^{\underline{j}}$$

 $a_j(G) = #\{$ partitions of V(G) into j independent subsets $\},$

$$P(G; k) = \sum_{1 \le j \le |V(G)|} (-1)^{j} b_{j}(G) k^{|V(G)| - j}$$

 $b_j(G) = \#\{j \text{-subsets of } E(G) \text{ containing no broken cycle}\}.$

 $uv \in E(G), \quad P(G; k) = P(G \setminus uv; k) - P(G/uv; k)$

Graph polynomials Graph homomorphisms

Independence polynomial

Definition by coefficients

$$I(G; x) = \sum_{1 \le j \le |V(G)|} b_j(G) x^j,$$

 $b_j(G) = #\{$ independent subsets of V(G) of size $j\}$.

 $v \in V(G), \quad I(G;x) = I(G - v;x) + xI(G - N[v];x)$

I(L(G); x) = matching polynomial of G

(Chudnovsky & Seymour, 2006) $K_{1,3} \not\subseteq_i G \Rightarrow I(G; x)$ real roots $b_j^2 \ge b_{j-1}b_{j+1}$, (implies $b_1, \ldots, b_{|V(G)|}$ unimodal) What are we looking at?

Sequences giving graph polynomials Building strongly polynomial graph sequences Interpretation schemes Open problems

Graph polynomials Graph homomorphisms

Definition

Graphs G, H. $f: V(G) \rightarrow V(H)$ is a homomorphism from G to H if $uv \in E(G) \Rightarrow f(u)f(v) \in E(H)$.

Definition

H with adjacency matrix $(a_{s,t})$, weight $a_{s,t}$ on $st \in E(H)$,

$$\hom(G,H) = \sum_{f:V(G)\to V(H)} \prod_{uv\in E(G)} a_{f(u),f(v)}.$$

 $\begin{aligned} \hom(G,H) &= \#\{\text{homomorphisms from } G \text{ to } H\} \\ &= \#\{H\text{-colourings of } G\} \end{aligned}$

when H simple $(a_{s,t} \in \{0,1\})$ or multigraph $(a_{s,t} \in \mathbb{N})$

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The main question

Which sequences $(H_{k,\ell,...})$ of simple graphs are such that, for all graphs G, for each $k, \ell, \cdots \in \mathbb{N}$ we have

$$\hom(G, H_{k,\ell,\ldots}) = p(G; k, \ell, \ldots)$$

for polynomial p(G)?

Characterizing simple graph sequences $(H_{k,\ell,...})$ with this property gives straightforward characterization for multigraph sequences too (allowing multiple edges & loops).

What are we looking at?

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Graph polynomials Graph homomorphisms



Examples

Strongly polynomial sequences of graphs Counting induced subgraphs

Example 1



$\hom(G, K_k) = P(G; k)$

chromatic polynomial

Examples

Strongly polynomial sequences of graphs Counting induced subgraphs

Example 2



 (\mathcal{K}^1_k) hom $(\mathcal{G},\mathcal{K}^1_k)=k^{|V(\mathcal{G})|}$

Examples

Strongly polynomial sequences of graph Counting induced subgraphs

Example 3



Examples

Strongly polynomial sequences of graph Counting induced subgraphs

Example 4



Examples Strongly polynomial sequences of g Counting induced subgraphs

Example 5



 $(K_2^{\Box k}) = (Q_k)$ (hypercubes)

Proposition (Garijo, G., Nešetřil, 2013+)

 $hom(G, Q_k) = p(G; k, 2^k)$ for bivariate polynomial p(G)

Examples Strongly polynomial sequences of graphs Counting induced subgraphs

Definition

 (H_k) is strongly polynomial (in k) if $\forall G \exists$ polynomial p(G) such that $\hom(G, H_k) = p(G; k)$ for all $k \in \mathbb{N}$.

Since $\hom(G_1 \cup G_2, H) = \hom(G_1, H) \hom(G_2, H)$, suffices to consider *connected* G.

Example

- (K_k) , (K_k^1) . $(\overline{kK_2})$ are strongly polynomial
- (K_k^{ℓ}) is strongly polynomial (in k, ℓ)
- (Q_k) not strongly polynomial (but polynomial in k and 2^k)
- (C_k) , (P_k) not strongly polynomial (but eventually polynomial in k)

Examples Strongly polynomial sequences of graphs Counting induced subgraphs

Subgraph criterion for strongly polynomial

$$H_k ext{ simple:} ext{ hom}(G, H_k) = \sum_{\substack{S \subseteq_i H_k \ |V(S)| \leq |V(G)|}} ext{surv}(G, S)$$

 $= \sum_{S/\cong} \operatorname{sur}_{\mathsf{v}}(G,S) \ \#\{ \text{induced copies of } S \text{ in } H_k \}$

Proposition (de la Harpe & Jaeger 1995)

(H_k) is strongly polynomial \iff \forall connected $S \ \#$ {induced subgraphs $\cong S$ in H_k } polynomial in k

Examples Strongly polynomial sequences of graphs Counting induced subgraphs

Subgraph criterion for strongly polynomial

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$$= \sum_{S/\cong} \operatorname{sur}_{\mathsf{v}}(G, S) \, \#\{ \text{induced copies of } S \text{ in } H_k \}$$

(for each S want this polynomial in k)

Proposition (de la Harpe & Jaeger 1995)

(H_k) is strongly polynomial \iff \forall connected $S \ \#$ {induced subgraphs $\cong S$ in H_k } polynomial in k

Examples Strongly polynomial sequences of graphs Counting induced subgraphs

Example: chromatic polynomial



Examples Strongly polynomial sequences of graphs Counting induced subgraphs

Eventually polynomial but not strongly polynomial



$\hom(G, C_k) = \sum_{1 \leq j \leq \min\{|V(G)|, k-1\}} \operatorname{sur}_{V}(G, P_j) k + \operatorname{sur}_{V}(G, C_k)$

 $hom(C_3, C_3) = 6$, $hom(C_3, C_k) = 0$ when k = 2 or $k \ge 4$

Constructions

Loose threads up until a few months ago...



Constructions Loose threads up until a few months ago...

Proposition (de la Harpe & Jaeger, 1995; Garijo, G., Nešetřil, 2013+)

If (H_k) is strongly polynomial and H_k simple, then

- $(\overline{H_k})$ (complements),
- $(L(H_k))$ (line graphs),

are also strongly polynomial. Also, (ℓH_k) is strongly polynomial (in k and ℓ).

Proposition (Garijo, G., Nešetřil, 2013+)

If (H_k) is strongly polynomial, at most one loop each vertex of H_k , then

- (H_k^0) (remove all loops)
- (H_k^1) (add loops to make 1 loop each vertex)

are also strongly polynomial. More generally, (H_k^{ℓ}) is strongly polynomial (in k and ℓ).

Constructions

Loose threads up until a few months ago...

Proposition

If (F_j) , (H_k) are strongly polynomial, then

- $(F_j \cup H_k)$ (disjoint union)
- $(F_j + H_k)$ (join)
- $(F_j \times H_k)$ (direct/categorical product)
- (*F_j*[*H_k*]) (lexicographic product)

are strongly polynomial (in j and k).

Constructions

Loose threads up until a few months ago...

Example

Beginning with trivially strongly polynomial sequence (K_1) , following are also strongly polynomial:

- multiple: $(kK_1) = (\overline{K_k})$
- complement: (K_k) (chromatic polynomial)
- loop-addition: (K_k^{ℓ}) (Tutte polynomial)

• join: $(K_{k-j}^1 + K_j^\ell)$ (Averbouch–Godlin–Makowsky polynomial – includes Tutte polynomial, satisfies three-term recurrence in $\langle uv, /uv | and -u - v \rangle$

Constructions Loose threads up until a few months ago...

Question

Strongly polynomial sequences:

$$\blacktriangleright (\overline{K_j} + \overline{K_k}) = (K_{j,k})$$

•
$$(L(K_{j,k})) = (K_j \Box K_k)$$
 (Rook's graph)

 (F_j) , (H_k) strongly polynomial $\Rightarrow (F_j \Box H_k)$ strongly polynomial?

Constructions Loose threads up until a few months ago...

Definition

Generalized Johnson graph $J_{k,\ell,D}$, $D \subseteq \{0, 1, \dots, \ell\}$ vertices $\binom{[k]}{\ell}$, edge uv when $|u \cap v| \in D$

- Johnson graphs $D = \{k 1\}$
- Kneser graphs $D = \{0\}$

Proposition (de la Harpe & Jaeger, 1995; Garijo, G., Nešetřil, 2013+)

For every ℓ , D, sequence $(J_{k,\ell,D})$ is strongly polynomial (in k).

Question

Can generalized Johnson graphs be generated from simpler sequences by any of the constructions described in de la Harpe & Jaeger (1995) and Garijo, Goodall & Nešetřil (2013+)?

Constructions Loose threads up until a few months ago...



Recap Relational structures Example interpretations Everything?

Simple graph sequence (H_k) strongly polynomial iff

- $\forall G \exists$ polynomial $p(G) \forall k \in \mathbb{N}$ hom $(G, H_k) = p(G; k)$
- $\forall F \exists$ polynomial $q(F) \forall k \in \mathbb{N}$ ind $(F, H_k) = q(F; k)$

Unary operations \sim and binary operations * such that if simple graph sequences (F_j) and (H_k) are strongly polynomial then

- (\widetilde{H}_k) is strongly polynomial (e.g. complement, line graph)
- (*F_j* * *H_k*) is strongly polynomial in *j*, *k* (e.g. join, lexicographic product)

Recap Relational structures Example interpretations Everything?

Satisfaction sets

Quantifier-free formula ϕ with *n* free variables ($\phi \in QF_n$) with symbols from relational structure **H** with domain $V(\mathbf{H})$.

Satisfaction set
$$\phi(\mathbf{H}) = \{(v_1, \dots, v_n) \in V(\mathbf{H})^n : \mathbf{H} \models \phi\}.$$

e.g. for graph structure H (symmetric binary relation $x \sim y$ interpreted as x adjacent to y), and given graph G on n vertices,

$$\phi = \phi_G = \bigwedge_{ij \in E(G)} (v_i \sim v_j)$$

 $\phi_G(H) = \{ (v_1, \dots, v_n) : i \mapsto v_i \text{ is a homomorphism } G \to H \}$

 $|\phi_{G}(H)| = \hom(G, H).$

Recap Relational structures Example interpretations Everything?

Strongly polynomial sequences of structures

Definition

Sequence (\mathbf{H}_k) of relational structures strongly polynomial iff $\forall \phi \in QF \exists$ polynomial $r(\phi) \forall k \in \mathbb{N} |\phi(\mathbf{H}_k)| = r(\phi; k)$

Lemma

Equivalently,

- $\forall \mathbf{G} \exists \text{ polynomial } p(\mathbf{G}) \forall k \in \mathbb{N} \quad \hom(\mathbf{G}, \mathbf{H}_k) = p(\mathbf{G}; k), \text{ or }$
- $\forall \mathbf{F} \exists \text{ polynomial } q(\mathbf{F}) \forall k \in \mathbb{N} \quad ind(\mathbf{F}, \mathbf{H}_k) = q(\mathbf{F}; k).$

Transitive tournaments (\vec{T}_k) strongly polynomial sequence of digraphs (e.g. count induced substructures).

Recap Relational structures Example interpretations Everything?

Graphical QF interpretation schemes

I: Relational σ -structures **A** \longrightarrow Graphs H

Lemma

There is

$$\widetilde{\mathit{I}}: \phi \in \mathrm{QF}(\mathrm{Graphs}) \quad \longmapsto \quad \widetilde{\mathit{I}}(\phi) \in \mathrm{QF}(\sigma extsf{-structures})$$

such that

$$\phi(I(\mathbf{A})) = \widetilde{I}(\phi)(\mathbf{A})$$

In particular, (\mathbf{A}_k) strongly polynomial \Rightarrow $(H_k) = (I(\mathbf{A}_k))$ strongly polynomial.

Recap Relational structures Example interpretations Everything?

From graphs to graphs

- All previous constructions (complementation, line graph, disjoint union, join, direct product,...) special cases of interpretation schemes *I* from Marked Graphs (added unary relations) to Graphs.
- Cartesian product and other more complicated graph products are special kinds of such interpretation schemes too.

Recap Relational structures Example interpretations Everything?

- Generalized Johnson graphs (J_{k,ℓ,D}) arise as QF interpretations of transitive tournaments T
 [¯]_k
- Half-graphs are QF interpretations of a transitive tournament together with "marks" (unary relations used to specfiy "upper" + "lower" vertices) and so form a strongly polynomial sequence.



Recap Relational structures Example interpretations Everything?

• Intersection graphs of chords of a *k*-gon form a strongly polynomial sequence



(a) Square



(b) Pentagon



Recap Relational structures Example interpretations Everything?

Conjecture

All strongly polynomial sequences of graphs (H_k) can be obtained by QF interpretation of a "basic sequence" (disjoint union of marked transitive tournaments of size polynomial in k).

Recap Relational structures Example interpretations Everything?



Paley graphs Further problems

Prime power $q = p^d \equiv 1 \pmod{4}$ Paley graph $P_q = \text{Cayley}(\mathbb{F}_q, \text{non-zero squares})$, Quasi-random graphs: $\hom(G, P_q) / \hom(G, G_{q, \frac{1}{2}}) \to 1 \text{ as } q \to \infty$.

Proposition (Corollary to result of de la Harpe & Jaeger, 1995)

hom (G, P_q) is polynomial in q for series-parallel G. e.g. hom $(K_3, P_q) = \frac{q(q-1)(q-5)}{8}$

Prime $q \equiv 1 \pmod{4}$, $q = 4x^2 + y^2$, [Evans, Pulham, Sheehan, 1981]: $\hom(\mathcal{K}_4, P_q) = \frac{q(q-1)}{1536} \left((q-9)^2 - 4x^2 \right)$

Is $hom(G, P_q)$ polynomial in q and x for all graphs G?

Theorem (G., Nešetřil, Ossona de Mendez , 2014+)

If (H_k) is strongly polynomial then there are only finitely many terms belonging to a quasi-random sequence of graphs.

Paley graphs Further problems

- When is (Cayley(A_k, B_k)) polynomial in |A_k|, |B_k|, where B_k = -B_k ⊆ A_k?
 e.g. For D ⊂ N, sequence (Cayley(Z_k, ±D)) is polynomial iff D is finite or cofinite. (de la Harpe & Jaeger, 1995)
- ► Can (*H_k*) be verified to be strongly polynomial by testing hom(*G*, *H_k*) for *G* only in a restricted class of graphs? (yes, for connected graphs – but for a smaller class?)
- Which graph polynomials defined by strongly polynomial sequences of graphs satisfy a reduction formula (size-decreasing recurrence) like the chromatic polynomial and independence polynomial?
- Develop similar theory for hom(H_k, G) (e.g. hom(C_k, G) = ∑ λ^k, λ eigenvalues of G, determines characteristic polynomial of G by its roots).

Paley graphs Further problems

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Paley graphs Further problems

Three papers

 P. de la Harpe and F. Jaeger, Chromatic invariants for finite graphs: theme and polynomial variations, *Lin. Algebra Appl.* 226–228 (1995), 687–722

Defining graphs invariants from counting graph homomorphisms. Examples. Basic constructions.

- D. Garijo, A. Goodall, J. Nešetřil, Polynomial graph invariants from homomorphism numbers. 40pp. arXiv: 1308.3999 [math.CO]
 Further examples. New construction using tree representations of graphs.
- A. Goodall, J. Nešetřil, P. Ossona de Mendez, Strongly polynomial sequences as interpretation of trivial structures. 17pp. Preprint. General relational structures: counting satisfying assignments for quantifier-free formulas. Building new polynomial invariants by interpretation of "trivial" sequences of marked tournaments.